The current topic: Scheme	Announcements
 Introduction Object-oriented programming: Python Features, variables, numbers, strings, Booleans, while loops If statements, sequences, functions, modules Dictionaries, command-line arguments, files, classes, inheritance, polymorphism Exceptions, operator overloading, privacy Multiple inheritance, parameters and arguments, list comprehensions Regular expressions, doc strings Functional programming: Scheme Next up: Introduction Types and values Syntax and semantics Exceptions Logic programming: Prolog 	 Reminder: Lab 1 was due at 10:30am today. Late penalty is 20% per day. Not accepted after 10:30am on Wednesday. Reminder: Term Test 1 is on October 6th in GB405, not in the regular lecture room. You're allowed to have one double-sided aid sheet for the test. You must use standard letter-sized (that is, 8.5" x 11") paper. The aid sheet can be produced however you like (typed or handwritten). Bring your TCard. What's covered? Everything from the first seven lectures (that is, everything up to and including September 26th). Lab 1. An old Term Test 1 has been posted. The exercises at the end of each lecture are also good practice.
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Basic functional-programming concepts

- Reference: Sebesta, chapter 15.
- <u>Referential transparency</u>: The value of a function application is independent of the context in which it occurs.
 - The value of f(a, b, c) depends only on the values of f, a, b and c.
 - It does not depend on the global state of computation.
 - That is, *the point in time* at which the call is made has no effect on the value that is returned.
 - This makes it easier to understand what's happening in a program.
- All variables that appear inside a function must be parameters.
 - No variables can be declared or first introduced within a function body.
 - A function has no access to global variables.

Basic functional-programming concepts

- *No assignment*: The concept of assigning a new value to a variable is not part of functional programming.
 - There are no explicit assignment statements.
 - Variables are bound to values only through the association of arguments to parameters in function calls.
 - Function calls have no side effects.
 - They do not alter the value of their arguments, or make any other changes to the global state.
- Conclusion: There is no need to consider global state.

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Control flow		Functio	onal programming: An example		
		 Recall how we comp 	outed the n-th Fibonacci number in Python:		
 Control is through function calls and conditional expressions only. There are no loops of any kind. 		<pre>def fib(n): # We'll assume n >= 1.</pre>			
Recursion is used to achieve repetition.		<pre>prev = 1 # cur = 1 for i in ran prev, cur return cur • To re-write this funct need to get rid of the - How can we do this?</pre>	<pre>first Fibonacci number # second Fibonacci number .ge(2, n): = cur, prev+cur tion using a functional programming approach, we e assignment statements and the loop. ?</pre>		
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Functional programming: An example Storage management • The same function, using a functional programming approach (but still in • All variables are function parameters, so structures that have a lifetime Python): beyond a function call can't be referred to. - A structure returned by a function is never assigned to a variable, so there's no way to explicitly free it. def fib(n): if n <= 2: - e.g. Suppose (for the sake of this discussion) that in the example on the previous slide, fib() returned a list. How would we free the lists returned by fib(n-1) return 1 and fib(n-2)? else: return fib(n-1) + fib(n-2) • Garbage collection is needed to dispose of such structures. Observe that the functional version uses an if statement and recursion. Fall 2008 Scheme: Introduction 8

Functions as first-class values Functions as "values" in C Recall that first-class "objects" are things that a language allows • C allows *pointers* to functions, like this: programs to : double integrate(double (*f)(double x), double a, double b) { - declare ... sum += (*f)(a + k*step); ... - assign values to } - return as function values - return as the value of an expression Function pointers are often used in tables to associate an action with an - pass as an argument index. - store in a data structure - create without a name • So, C function *pointers* can be stored in a data structure as a value. For example, integers are first-class values in most programming But in C, unnamed functions cannot exist as values, though there can languages. be unnamed function pointers. In functional programming languages, functions are first-class values. Functions are **not** first-class values in C. - Aside: Python also mostly treats functions as first-class values. - And they're not first-class values in Java either, which doesn't even have function • Why "mostly"? There are some restrictions on creating unnamed functions. pointers. Fall 2008 10 Fall 2008 Scheme: Introduction 9 Scheme: Introduction Common LISP **History: LISP** "LISt Processing" LISP implementations began to stray - from the original semantics • A functional language developed by John McCarthy in the late 1950s for use in Al. - from each other Semantics were redefined in Common LISP. Semantics (meaning) based on λ-calculus.

- The $\lambda\mbox{-}calculus$ is a theoretical model of computation
- Functions operate on lists or atomic symbols.
 - Programs consist of "S-expressions".
 - five basic functions: cons, car, cdr, equal, atom
 - one conditional construct: cond
- Programs and data are both S-expressions.

- Common LISP is now the standard version of LISP. – Designed by committee in 1980s.
 - -
- Simple syntax, but a large language.

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Scheme	The Scheme we'll use		
Created in 1975 by Gerald Sussman and Guy Steele.	 There are many flavours of Scheme, with quite a bit of variation among them. But not as much variation as among the LISPs. 		
 A version of LISP Simple syntax, small language Close to the initial semantics of LISP Provides basic list-processing tools Allows functions to be first-class objects Not the same as LISP! 	 We'll use MzScheme. Use the command "mzscheme" to launch the MzScheme interpreter on ECF. Version 202 is installed on ECF. This is not the most recent version. DrScheme is an IDE that works with MzScheme. Use the command "drscheme" on ECF. Make sure the selected language is MzScheme. 		
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The Scheme we'll use

- DrScheme can be downloaded from
 <u>http://download.plt-scheme.org/drscheme/v202.html</u>
 - This download also includes MzScheme.
 - Note that the above link is for version 202, the version we're using on ECF.
- Your assignments must run using MzScheme as installed on ECF.

Scheme references

- Sebesta, Section 15.5.
- Dybvig, The Scheme Programming Language, 2003.
 - This textbook can be read online at: <u>http://www.scheme.com/tspl3/</u>
- "Teach Yourself Scheme in Fixnum Days", a tutorial at: <u>http://www.ccs.neu.edu/home/dorai/t-y-scheme/t-y-scheme.html</u>

A Scheme program The basic structure of a Scheme (or LISP) program A program consists of A program consists of S-expressions ("symbolic expressions") written in parenthesized prefix form. - a set of function definitions - "Prefix form": The name of the function appears before the arguments to the - an expression to be evaluated function, even for mathematical functions like +, -, *, and /. - The general form of an S-expression in prefix-form is: • A simple example that defines function abs-val and then calls this '(' <function> <arg1> <arg2> ... <argn> ')' function (note that ">" is the interpreter's prompt): > (define (abs-val x) • Examples of S-expressions in prefix-form: (if (>= x 0))х (+ 4 5) ; A simple example. (-x))• Note that ";" denotes a comment in Scheme. > (abs-val (-35))(+(*345)(-53))2 17 Fall 2008 Scheme: Introduction Fall 2008 Scheme: Introduction 18

S-expressions

- To evaluate an expression:
 - Evaluate the <function> to a value that is a function.
 - Evaluate each <argi> to obtain its value.
 - Apply the function value to these argument values.
- Evaluating our examples:
 - (+ 4 5)
 - 4 and 5 evaluate to themselves.
 - \bullet Applying "+" to 4 and 5 results (not surprisingly) in the value 9.

(+ (* 3 4 5) (- 5 3))

- Evaluate (* 3 4 5) to obtain the value 60.
- Evaluate (- 5 3) to obtain value 2.
- Evaluate (+ 60 2) to obtain value 62.

S-expressions

- So far we've seen S-expressions that are lists (that is, enclosed in brackets).
- Atoms are also S-expressions.
- An atom is anything that is not a non-empty list. Atoms include: - Numbers.
 - Strings.
 - Symbols.
 - The empty list.
 - Booleans.

S-expressions		Built-in functions	
 Examples: #t ; true #f ; false () ; the empty list (a b c) ; a list containing 3 symbols (a (b c) d) ; a nested list ((a b c) (d e (f))) (1 (b) 2) Observe that lists can be nested. 	• c: • c: • (c el • q • e • n	 ar : Returns the first element of a list. dr : Returns the rest of a list. cons <element> <list>) : Returns a new list formed by adding an lement to the front of the given list.</list></element> uote or ' : Produces constants. q? : Returns true iff given two identical atoms. ull? : Returns true iff given an empty list. 	
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car		cdr		
• car returns the first element of a list.		• cdr returns the rest of a list	(that is, everything except the first element).	
> (car '(a b c)) a		> (cdr '(a b c)) (b c)		
> (car '(1 2 3 4)) 1		<pre>> (cdr '(1 2 3 4)) (2 3 4)</pre>		
> (car '((a) b (c d))) (a)		> (cdr '((a) b (c d)) (b (c d)))	
> (car '((a b) c d)) (a b)		> (cdr '((a b) c d)) (c d)		
> (car '()) ; Causes an error.		> (cdr '(1)) ()		
		> (cdr '()) ; Causes	an error.	
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car and cdr cons • cons returns a new list formed by adding a given element to the front of a given list. • car and cdr can break up any list: > (cons 'a '()) > (car (cdr (cdr '((a) b (c d))))) (a) (c d) > (cons 'd '(e)) > (cdr '((a) b (c d))) (d e) (b (c d)) > (cons '(a b) '(c d)) > (cdr '(b (c d))) ((a b) c d) ((c d)) > (cons '(a b c) '((a) b)) > (car '((c d))) ((a b c) (a) b) (c d) > (cons (car '(a b c)) (cdr '(a b c))) > (caddr '((a) b (c d))); caddr = car cdr cdr (a b c) (c d) Scheme: Introduction 25 Fall 2008 Scheme: Introduction Fall 2008 26

Exercises

- Run mzscheme on ECF.
 - Try evaluating simple expressions like those covered in this lecture.
 - Experiment with car, cdr, and cons.
 - Try out the built-in combinations of car and cdr, like cdddr, caddr, etc.
 How many cars and cdrs can be combined into a single built-in function in mzscheme?
- Run drscheme on ECF.

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- Set the language to MzScheme.
- Repeat the above exercises, this time in $\ensuremath{\mathsf{DrScheme}}$.