

Let's make another
instance of our
framework!

Very Busy Expressions

if $[a > b]^1$ then $([x := b - a]^2; [y := a - b]^3)$ else $([y := b - a]^4; [x := a - b]^5)$



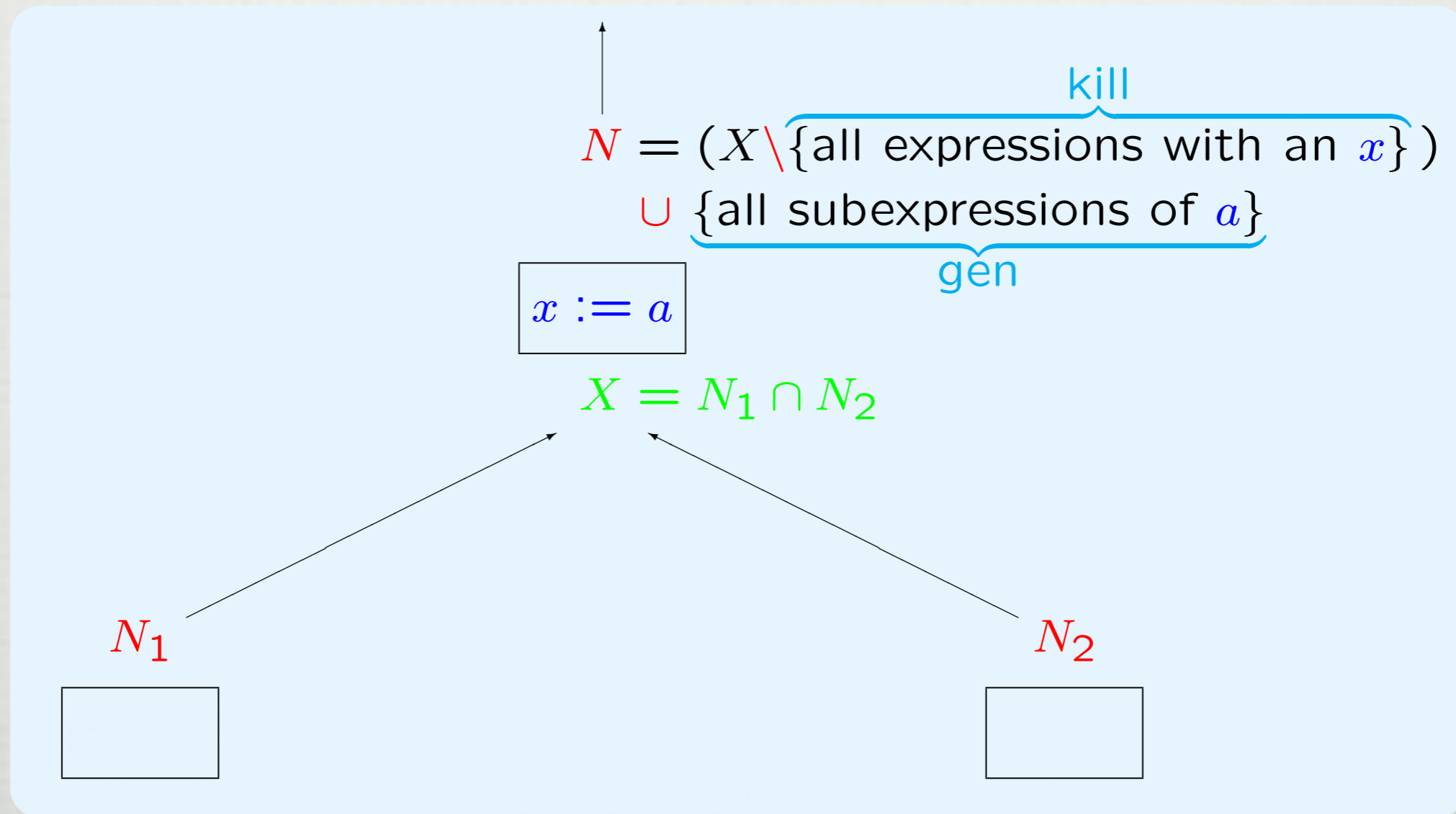
$[t1 := b - a]^A; [t2 := a - b]^B;$
if $[a > b]^1$ then $([x := t1]^2; [y := t2]^3)$ else $([y := t1]^4; [x := t2]^5)$

An expression is **very busy** at the **exit** from a label if, no matter what path is taken from the label, the expression is always used before any of the variables occurring in it are redefined.

Check List

- Define the semi-lattice: dataflow facts and how to combine them!
 - Decide on the direction of the analysis: forward vs backward.
 - Sanity check: the corresponding order should make sense!
 - Decide on the initial values.
- Design the transfer functions:
 - How does each statement affect the dataflow facts?
 - Sanity check: Monotonicity!

The Design Process



Very Busy Expressions: Formal Setup

Dataflow Facts: $D = \mathcal{P}(Exp)$

Domain: complete meet semi-lattice (D, \cap, D)

unlike **live variables**: here we want the greatest fixed point!

Direction: **Backward**

Transfer Functions:

$$VB_{\bullet}(l) = \begin{cases} \emptyset & l = \textit{exit} \\ \bigcap_{(l,l') \in \textit{flow}} VB_{\bullet}(l') & \textit{otherwise} \end{cases}$$

$$VB_{\circ}(l) = RD_{\bullet}(l) \setminus \{ \textit{exp} \mid \textit{var}(\textit{exp}) \cap \textit{write}(l) \neq \emptyset \} \\ \cup \textit{computed}(l)$$

Reaching Definitions

	Live Variables	Reaching Definitions	Very Busy Expressions
L	$\mathcal{P}(Var)$	$\mathcal{P}(Var \times Loc)$	$\mathcal{P}(Exp)$
\sqcap	\cup	\cup	\cap
\sqsubseteq	\supseteq	\supseteq	\subseteq
\perp	\emptyset	\emptyset	Exp
	Backward	Forward	Backward

Reaching Definitions

	←	$\{(x, ?), (y, ?), (z, ?)\}$
$[y := x]^1;$	←	$\{(x, ?), (y, 1), (z, ?)\}$
$[z := 1]^2;$	←	$\{(x, ?), (y, 1), (y, 5), (z, 2), (z, 4)\}$
while $[y > 0]^3$ do	←	$\{(x, ?), (y, 1), (y, 5), (z, 2), (z, 4)\}$
$[z := z * y]^4;$	←	$\{(x, ?), (y, 1), (y, 5), (z, 4)\}$
$[y := y - 1]^5$	←	$\{(x, ?), (y, 5), (z, 4)\}$
od;	←	$\{(x, ?), (y, 1), (y, 5), (z, 2), (z, 4)\}$
$[y := 0]^6$	←	$\{(x, ?), (y, 6), (z, 2), (z, 4)\}$

Reaching Definitions: Formal Setup

Dataflow Facts: $D = Vars \times Locs$

Domain: (D, \cup, \emptyset, D)

Direction: **Forward**

Transfer Functions:

$$RD_{\circ}(l) = \begin{cases} Var \times \{?\} & l = init \\ \bigcup_{(l',l) \in flow} RD_{\bullet}(l') & \text{otherwise} \end{cases}$$

$$RD_{\bullet}(l) = RD_{\circ}(l) \setminus write(l) \times Locs \cup write(l) \times \{l\}$$