

144 Let x be an integer variable. Prove the refinement

- (a) $x'=0 \iff \text{if } x=0 \text{ then } ok \text{ else } x:=x-1. x'=0 \text{ fi}$
- (b) $P \iff \text{if } x=0 \text{ then } ok \text{ else } x:=x-1. t:=t+1. P \text{ fi}$
where $P = x'=0 \wedge \text{if } x \geq 0 \text{ then } t' = t+x \text{ else } t' = \infty \text{ fi}$

After trying the question, scroll down to the solution.

$$(a) \quad x' = 0 \iff \text{if } x = 0 \text{ then } ok \text{ else } x := x - 1. \quad x' = 0 \text{ fi}$$

§ By Cases. First case:

$$\begin{aligned} & x = 0 \wedge ok \Rightarrow x' = 0 && \text{expand } ok \\ = & x = 0 \wedge x' = x \wedge t' = t \Rightarrow x' = 0 && \text{context} \\ = & x = 0 \wedge x' = x \wedge t' = t \Rightarrow x = x && \text{reflexive} \\ = & x = 0 \wedge x' = x \wedge t' = t \Rightarrow \top && \text{base} \\ = & \top && \end{aligned}$$

Last case:

$$\begin{aligned} & x \neq 0 \wedge (x := x - 1. \quad t := t + 1. \quad x' = 0) && \text{substitution twice} \\ = & x \neq 0 \wedge x' = 0 && \text{specialization} \\ \Rightarrow & x' = 0 && \end{aligned}$$

$$(b) \quad P \iff \text{if } x = 0 \text{ then } ok \text{ else } x := x - 1. \quad t := t + 1. \quad P \text{ fi}$$

where $P = x' = 0 \wedge \text{if } x \geq 0 \text{ then } t' = t + x \text{ else } t' = \infty \text{ fi}$

§ By Parts, and part (a) proved one part, so we just have to prove

$$\text{if } x \geq 0 \text{ then } t' = t + x \text{ else } t' = \infty \text{ fi} \iff \text{if } x = 0 \text{ then } ok \text{ else } x := x - 1. \quad t := t + 1. \quad P \text{ fi}$$

And we prove it by Cases. First case:

$$\begin{aligned} & x = 0 \wedge ok \Rightarrow \text{if } x \geq 0 \text{ then } t' = t + x \text{ else } t' = \infty \text{ fi} && \text{expand } ok \\ = & x = 0 \wedge x' = x \wedge t' = t \Rightarrow \text{if } x \geq 0 \text{ then } t' = t + x \text{ else } t' = \infty \text{ fi} && \text{context} \\ = & x = 0 \wedge x' = x \wedge t' = t \Rightarrow \text{if } 0 \geq 0 \text{ then } t = t + 0 \text{ else } t' = \infty \text{ fi} && \text{reflexive, case base} \\ = & x = 0 \wedge x' = x \wedge t' = t \Rightarrow t = t && \text{reflexive} \\ = & x = 0 \wedge x' = x \wedge t' = t \Rightarrow \top && \text{base} \\ = & \top && \end{aligned}$$

Last case:

$$\begin{aligned} & x \neq 0 \wedge (x := x - 1. \quad t := t + 1. \quad \text{if } x \geq 0 \text{ then } t' = t + x \text{ else } t' = \infty \text{ fi}) && \text{substitution twice} \\ = & x \neq 0 \wedge \text{if } x - 1 \geq 0 \text{ then } t' = t + 1 + x - 1 \text{ else } t' = \infty \text{ fi} && \text{simplify} \\ = & x \neq 0 \wedge \text{if } x > 0 \text{ then } t' = t + x \text{ else } t' = \infty \text{ fi} && \text{context: } x \neq 0 \Rightarrow (x > 0 \equiv x \geq 0) \\ = & x \neq 0 \wedge \text{if } x \geq 0 \text{ then } t' = t + x \text{ else } t' = \infty \text{ fi} && \text{specialization} \\ \Rightarrow & \text{if } x \geq 0 \text{ then } t' = t + x \text{ else } t' = \infty \text{ fi} && \end{aligned}$$