

274 (smallest common item) Given two sorted lists having at least one item in common, write a program to find the smallest item occurring in both lists.

After trying the question, scroll down to the solution.

§ Let the two lists be A and B , and let a and b be natural variables to act as indexes of the two lists. The specification is R , defined as

$$\begin{aligned} R = & \quad A(0,..a') \wedge B(0,..b') = null \\ & \quad \wedge A a' = B b \\ & \quad \wedge t' = t + a' + b' \end{aligned}$$

which says there isn't a common item before indexes a' and b' , and there is one at those indexes, and the sum of those indexes is the execution time. We also need the specification Q , defined as

$$\begin{aligned} Q = & \quad A(a,..a') \wedge B(b,..b') = null \\ & \quad \wedge A a' = B b' \\ & \quad \wedge t' = t + a' - a + b' - b \end{aligned}$$

which says the same thing, but starting at indexes a and b . The refinements are

$$\begin{aligned} R & \leftarrow a := 0. b := 0. Q \\ Q & \leftarrow \text{if } A a < B b \text{ then } a := a + 1. t := t + 1. Q \\ & \quad \text{else if } B b < C c \text{ then } b := b + 1. t := t + 1. Q \\ & \quad \text{else ok fi fi} \end{aligned}$$

The proof of the first refinement is immediate after using the Substitution Law 2 times. The proof of the last refinement breaks into 9 pieces (3 conjuncts \times 3 cases). Let's start with the first conjunct of Q and the first case of the **if**.

$$\begin{aligned} & A a < B b \wedge (a := a + 1. t := t + 1. A(a,..a') \wedge B(b,..b') = null) \quad \text{substitution law twice} \\ = & \quad A a < B b \wedge A(a + 1,..a') \wedge B(b,..b') = null \end{aligned}$$

Because B is sorted, $B b$ is the minimum item of $B(b,..b')$.

And since $A a < B b$, therefore $A a$ is unequal to any item in $B(b,..b')$.

$$\Rightarrow A(a,..a') \wedge B(b,..b') = null$$

Now we prove the middle conjunct of Q with the same first case of the **if**.

$$\begin{aligned} & A a < B b \wedge (a := a + 1. t := t + 1. A a' = B b') \quad \text{substitution law twice} \\ = & \quad A a < B b \wedge A a' = B b' \quad \text{specialize} \\ \Rightarrow & \quad A a' = B b' \end{aligned}$$

Now we prove the last conjunct of Q with the same first case of the **if**.

$$\begin{aligned} & A a < B b \wedge (a := a + 1. t := t + 1. t' = t + a' - a + b' - b) \quad \text{substitution law twice} \\ = & \quad A a < B b \wedge t' = t + 1 + a' - a - 1 + b' - b \quad \text{simplify and specialize} \\ \Rightarrow & \quad t' = t + a' - a + b' - b \end{aligned}$$

The proof of the middle case is exactly the same. That leaves the last case.

$$\begin{aligned} & A a \geq B b \wedge B b \geq A a \wedge ok \quad \text{antisymmetry, and expand } ok \\ & A a = B b \wedge a' = a \wedge b' = b \wedge t' = t \\ \Rightarrow & \quad Q \end{aligned}$$