

290 (factorial space) We can compute $x := n!$ (factorial) as follows.

$x := n! \iff \mathbf{if } n=0 \mathbf{ then } x := 1 \mathbf{ else } n := n-1. x := n!. n := n+1. x := x \times n \mathbf{ fi}$

Each call $x := n!$ pushes a return address onto a stack, and each return pops an address from the stack. Add a space variable s and a maximum space variable m , with appropriate assignments to them in the program. Find and prove an upper bound on the maximum space used.

§ $m \geq s \Rightarrow m' = m \uparrow (s+n) \iff$
 $\mathbf{if } n=0 \mathbf{ then } x := 1$
 $\mathbf{else } n := n-1.$
 $s := s+1. m := m \uparrow s. m \geq s \Rightarrow m' = m \uparrow (s+n). s := s-1.$
 $n := n+1. x := x \times n \mathbf{ fi}$

Proof: by cases. First case:

$(n=0 \wedge (x:=1) \Rightarrow (m \geq s \Rightarrow m' = m \uparrow (s+n)))$ portation and expand
 $= n=0 \wedge x'=1 \wedge n'=n \wedge s'=s \wedge m'=m \wedge m \geq s \Rightarrow m' = m \uparrow (s+n)$ context
 $= n=0 \wedge x'=1 \wedge n'=n \wedge s'=s \wedge m'=m \wedge m \geq s \Rightarrow \top$ base
 $= \top$

Second case:

$(m \geq s \Rightarrow m' = m \uparrow (s+n))$
 $\iff (n \neq 0 \wedge (n := n-1. s := s+1.$
 $m := m \uparrow s. m \geq s \Rightarrow m' = m \uparrow (s+n). s := s-1.$
 $n := n+1. x := x \times n))$ portation
 $= n \neq 0 \wedge m \geq s$
 $\wedge (n := n-1. s := s+1. m := m \uparrow s. m \geq s \Rightarrow m' = m \uparrow (s+n).$
 $s := s-1. n := n+1. x := x \times n)$
 $\Rightarrow m' = m \uparrow (s+n)$
 three substitutions; expand final assignment and two more substitutions
 $= n \neq 0 \wedge m \geq s$
 $\wedge (m \uparrow (s+1) \geq s+1 \Rightarrow m' = m \uparrow (s+1) \uparrow (s+1+n-1).$
 $x' = x \times (n+1) \wedge n' = n+1 \wedge s' = s-1 \wedge m' = m)$
 $\Rightarrow m' = m \uparrow (s+n)$ simplify first \uparrow to \top and $+1-1$ to 0 .
 In the context where natural n is not 0 , $s+n \geq s+1$ so remove $s+1$ from \uparrow .
 $= n \neq 0 \wedge m \geq s$
 $\wedge (m' = m \uparrow (s+n). x' = x \times (n+1) \wedge n' = n+1 \wedge s' = s-1 \wedge m' = m)$
 $\Rightarrow m' = m \uparrow (s+n)$ Eliminate $.$ and then use one-point.
 $= n \neq 0 \wedge m \geq s \wedge m' = m \uparrow (s+n)$
 $\Rightarrow m' = m \uparrow (s+n)$ specialize
 $= \top$