

373 Prove the law  $nat = 0,..^\infty$  from the other laws.

After trying the question, scroll down to a solution attempt.

§ I'll start by proving half of it.

$nat: 0,..\infty$

In this half, I am trying to prove  $nat$  isn't too big, so I'll need  $nat$  induction

$0, B+1: B \Rightarrow nat: B$

Here we go:

$nat: 0,..\infty$	
$\Leftarrow 0, (0,..\infty)+1: 0,..\infty$	use induction with $0,..\infty$ for $B$
$= 0, 0+1,..\infty+1: 0,..\infty$	addition distributes over bunch union
$= 0, 1,..\infty: 0,..\infty$	arithmetic; $\infty$ absorbs additions
$= 0,..\infty: 0,..\infty$	reflexive
$= \top$	

The other half is

$0,..\infty: nat$

In this half, I am trying to prove that all of  $0,..\infty$  is in  $nat$ , so I should need  $nat$  construction.

$0, nat+1 : nat$

So maybe I should prove

$0,..\infty: 0, nat+1$

and then the result will follow by transitivity. Also, these laws look like they might be useful:

$x, y: xint \wedge x \leq y \Rightarrow (i: x,..y = i: xint \wedge x \leq i < y)$  interval law

$A: B = \forall x: A \cdot x: B$  inclusion law

$V: W = \forall v: V \cdot \exists w: W \cdot v=w$  bunch-element conversion law

$A: B \wedge B: C \Rightarrow A: C$  transitivity

But I am unable to find a proof.

Here's an attempt to prove both halves together.

$\top$	identity
$= \forall i: xint \cdot \top$	interval law
$= \forall i: xint \cdot x, y: xint \wedge x \leq y \Rightarrow (i: x,..y = i: xint \wedge x \leq i < y)$	specialize
$\Rightarrow \forall i: xint \cdot 0, \infty: xint \wedge 0 \leq \infty \Rightarrow (i: 0,..\infty = i: xint \wedge 0 \leq i < \infty)$	antecedent all $\top$
$= \forall i: xint \cdot (i: 0,..\infty = i: xint \wedge 0 \leq i < \infty)$	context provided by quantification
$(\text{quantifier applies to a function; function variable introduction is axiom in body})$	
$= \forall i: xint \cdot (i: 0,..\infty = \top \wedge 0 \leq i < \infty)$	identity
$= \forall i: xint \cdot (i: 0,..\infty = 0 \leq i < \infty)$	I can't justify this step
$= \forall i: xint \cdot (i: 0,..\infty = i: nat)$	I can't justify this step either
$= \forall i: xint \cdot 0,..\infty = nat$	idempotent
$= 0,..\infty = nat$	

You might just say it's obvious that  $nat = 0,..\infty$ , so why do we have to prove it? We have an application for our math: formal methods of software development. For that application, we want some binary expressions to be theorems (for example,  $nat = 0,..\infty$ ), and we want other binary expressions to be antitheorems (for example,  $nat \neq 0,..\infty$ ), and there are other binary expressions that we don't care about (for example,  $0/0 = 1$ ). If we say  $nat = 0,..\infty$  is obvious, that just means it's obvious that we want it to be a theorem. If it cannot be proven, we need to add it to the theory. I have added it, but it may be provable from the other laws.