

390 Let $A \setminus B$ be the difference between bunch A and bunch B . The operator \setminus has precedence level 4, and is defined by the axiom

$$x: A \setminus B = x: A \wedge \neg x: B$$

For each of the following fixed-point equations, what does recursive construction yield? Does it satisfy the fixed-point equation?

(a) $Q = nat \setminus (Q+3)$
 \S $Q_0 = null$
 $Q_1 = nat \setminus (null+3) = nat \setminus null = nat$
 $Q_2 = nat \setminus (nat+3) = 0, 1, 2$
 $Q_3 = nat \setminus ((0, 1, 2)+3) = nat \setminus (3, 4, 5) = 0, 1, 2, nat+6$
 $Q_4 = nat \setminus ((0, 1, 2, nat+6)+3) = nat \setminus (3, 4, 5, nat+9) = 0, 1, 2, 6, 7, 8$
 $Q_5 = nat \setminus ((0, 1, 2, 6, 7, 8)+3) = nat \setminus (3, 4, 5, 9, 10, 11)$
 $= 0, 1, 2, 6, 7, 8, nat+12$

Time for a guess. It looks like there are two patterns: the even index pattern and the odd index pattern. So I guess

$$Q_{2 \times n} = 6 \times (0, \dots, n) + (0, \dots, 3)$$

$$Q_{2 \times n + 1} = 6 \times (0, \dots, n) + (0, \dots, 3), (6 \times n, \dots, \infty)$$

From the even case, I propose

$$Q_\infty = 6 \times nat + (0, \dots, 3)$$

and now I have to check whether it satisfies the equation. Starting with the right side,

$$nat \setminus (Q_\infty + 3)$$

$$= nat \setminus (6 \times nat + (0, \dots, 3) + 3)$$

$$= nat \setminus (6 \times nat + (3, \dots, 6))$$

$$= nat \setminus ((0, 6, 12, 18, 24, \dots) + (3, \dots, 6))$$

$$= nat \setminus (3, 4, 5, 9, 10, 11, 15, 16, 17, 21, 22, 23, 27, 28, 29, \dots)$$

$$= 0, 1, 2, 6, 7, 8, 12, 13, 14, 18, 19, 20, 24, 25, 26, \dots$$

$$= (0, 6, 12, 18, 24, \dots) + (0, \dots, 3)$$

$$= 6 \times nat + (0, \dots, 3)$$

$$= Q_\infty$$

So it does satisfy the equation. From the odd case, we can't make a proposal because we can't simplify ∞, \dots, ∞ .

(b) $D = 0, (D+1) \setminus (D-1)$
 \S $D_0 = null$
 $D_1 = 0, (D_0+1) \setminus (D_0-1)$
 $= 0, (null+1) \setminus (null-1)$
 $= 0, null \setminus null$
 $= 0$
 $D_2 = 0, (D_1+1) \setminus (D_1-1)$
 $= 0, (0+1) \setminus (0-1)$
 $= 0, 1 \setminus -1$
 $= 0, 1$
 $D_3 = 0, (D_2+1) \setminus (D_2-1)$
 $= 0, ((0, 1)+1) \setminus ((0, 1)-1)$
 $= 0, (1, 2) \setminus (-1, 0)$
 $= 0, 1, 2$
 $D_4 = 0, (D_3+1) \setminus (D_3-1)$
 $= 0, ((0, 1, 2)+1) \setminus ((0, 1, 2)-1)$
 $= 0, (1, 2, 3) \setminus (-1, 0, 1)$
 $= 0, 2, 3$

$$\begin{aligned}
D_5 &= 0, (D_4+1) \setminus (D_4-1) \\
&= 0, ((0, 2, 3)+1) \setminus ((0, 2, 3)-1) \\
&= 0, (1, 3, 4) \setminus (-1, 1, 2) \\
&= 0, 3, 4 \\
D_6 &= 0, (D_5+1) \setminus (D_5-1) \\
&= 0, ((0, 3, 4)+1) \setminus ((0, 3, 4)-1) \\
&= 0, (1, 4, 5) \setminus (-1, 2, 3) \\
&= 0, 1, 4, 5 \\
D_7 &= 0, (D_6+1) \setminus (D_6-1) \\
&= 0, ((0, 1, 4, 5)+1) \setminus ((0, 1, 4, 5)-1) \\
&= 0, (1, 2, 5, 6) \setminus (-1, 0, 3, 4) \\
&= 0, 1, 2, 5, 6 \\
D_8 &= 0, (D_7+1) \setminus (D_7-1) \\
&= 0, ((0, 1, 2, 5, 6)+1) \setminus ((0, 1, 2, 5, 6)-1) \\
&= 0, (1, 2, 3, 6, 7) \setminus (-1, 0, 1, 4, 5) \\
&= 0, 2, 3, 6, 7
\end{aligned}$$

It's still hard to see the patterns, so maybe we have to go a bit farther. Then we see

$$\begin{aligned}
D_{4 \times n + 1} &= 0, 4 \times (0, \dots, n) + (3, 4) \\
D_{4 \times n + 2} &= 0, 1, 4 \times (0, \dots, n) + (4, 5) \\
D_{4 \times n + 3} &= 0, 1, 2, 4 \times (0, \dots, n) + (5, 6) \\
D_{4 \times n + 4} &= 0, 2, 3, 4 \times (0, \dots, n) + (6, 7)
\end{aligned}$$

We have a choice of four possible answers for D_∞ , but none of them satisfies the equation. Recursive construction fails.

(c)

$$E = nat \setminus (E+1)$$

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$$\begin{aligned}
E_0 &= null \\
E_1 &= nat \\
E_2 &= 0 \\
E_3 &= 0, nat+2 \\
E_4 &= 0, 2 \\
E_5 &= 0, 2, nat+4 \\
E_{2 \times n} &= 2 \times (0, \dots, n) \\
E_{2 \times n + 1} &= 2 \times (0, \dots, n), nat+2 \times n
\end{aligned}$$

From the even case, we propose

$$E_\infty = 2 \times nat$$

which satisfies the equation. From the odd case, we propose

$$E_\infty = 2 \times nat, \infty$$

which does not satisfy the equation.

(d)

$$F = 0, (nat \setminus F)+1$$

§

$$\begin{aligned}
F_0 &= null \\
F_1 &= nat \\
F_2 &= 0 \\
F_3 &= 0, nat+2 \\
F_4 &= 0, 2 \\
F_5 &= 0, 2, nat+4 \\
F_{2 \times n} &= 2 \times (0, \dots, n) \\
F_{2 \times n + 1} &= 2 \times (0, \dots, n), nat+2 \times n
\end{aligned}$$

From the even case, we propose

$$F_\infty = 2 \times nat$$

which satisfies the fixed-point equation. From the odd case, we propose

$$F_\infty = 2 \times \text{nat}, \infty$$

which does not satisfy the fixed-point equation.