- 417 Let the state consist of binary variables b and c. Let  $W = \mathbf{if} b \mathbf{then} P. W \mathbf{else} ok \mathbf{fi}$  $X = \mathbf{if} b \mathbf{v} c \mathbf{then} P. X \mathbf{else} ok \mathbf{fi}$
- (a) Find a counterexample to W. X = X.
- (b) Now let W and X be the weakest solutions of those equations, and prove W.X = X.

After trying the question, scroll down to the solution.

(a) Find a counterexample to W.X = X. Let P = ok, let  $\overline{W} = c := b \lor c$ , and let  $X = b := \bot$ . We have to check that the given § equations are satisfied. First, the W equation. if b then P. W else ok fi expand.  $\perp$  is identity for v = **if** b **then** ok.  $c := b \lor c$  **else**  $c := \bot \lor c$  **fi** ok is identity. In the else part, b is  $\perp$ . = if b then  $c := b \lor c$  else  $c := b \lor c$  fi case idempotent = WNext, check the X equation. if  $b \lor c$  then P. X else ok fi In the else part, b is  $\perp$ . =if  $b \lor c$  then ok.  $b := \bot$  else  $b := \bot$  fi ok is identity, case idempotent = X So both equations are satisfied. Now we check that  $W \cdot X = X$  is not satisfied. W.X=  $c := b \lor c$ .  $b := \bot$ expand final assignment =  $c := b \lor c$ .  $\neg b' \land c' = c$ substitution law =  $\neg b' \land c' = b \lor c$ And X =  $b := \bot$  $\neg b' \land c' = c$ \_ When the initial state is  $b \wedge \neg c$ , (W.X) leaves c' with value  $\top$ , but X leaves c' with value  $\perp$ .

(b) Now let W and X be the weakest solutions of those equations, and prove W.X = X.