

455 A theory provides three names: *set* , *flip* , and *ask* . It is presented by an implementation. Let $u: bin$ be the user's variable, and let $v: bin$ be the implementer's variable. The axioms are

$$set = v := \top$$

$$flip = v := \neg v$$

$$ask = u := v$$

- (a)✓ Replace v with $w: nat$ according to the data transformer $v = even\ w$.
- (b) Replace v with $w: nat$ according to the data transformer $(w=0 \Rightarrow v) \wedge (w=1 \Rightarrow \neg v)$. Is anything wrong?
- (c) Replace v with $w: nat$ according to $(v \Rightarrow w=0) \wedge (\neg v \Rightarrow w=1)$. Is anything wrong?

After trying the question, scroll down to the solution.

(a) \checkmark Replace v with $w: \text{nat}$ according to the data transformer $v = \text{even } w$.
 \S see book Section 7.2

(b) Replace v with $w: \text{nat}$ according to the data transformer $(w=0 \Rightarrow v) \wedge (w=1 \Rightarrow \neg v)$. Is anything wrong?

\S Operation *set* becomes

$$\begin{aligned} & \forall v. (w=0 \Rightarrow v) \wedge (w=1 \Rightarrow \neg v) \Rightarrow \exists v'. (w'=0 \Rightarrow v') \wedge (w'=1 \Rightarrow \neg v') \wedge (v := \top) \\ & = u'=u \wedge w' \neq 1 \end{aligned}$$

Operation *flip* becomes

$$\begin{aligned} & \forall v. (w=0 \Rightarrow v) \wedge (w=1 \Rightarrow \neg v) \Rightarrow \exists v'. (w'=0 \Rightarrow v') \wedge (w'=1 \Rightarrow \neg v') \wedge (v := \neg v) \\ & = u'=u \wedge (w \neq 0 \Rightarrow w' \neq 1) \wedge (w \neq 1 \Rightarrow w' \neq 0) \end{aligned}$$

Operation *ask* becomes

$$\begin{aligned} & \forall v. (w=0 \Rightarrow v) \wedge (w=1 \Rightarrow \neg v) \Rightarrow \exists v'. (w'=0 \Rightarrow v') \wedge (w'=1 \Rightarrow \neg v') \wedge (u := v) \\ & = (w \neq 0 \Rightarrow w' \neq 0 \wedge \neg u') \wedge (w \neq 1 \Rightarrow w' \neq 1 \wedge u') \\ & = (w=0 \wedge w' \neq 1 \wedge \neg u') \vee (w=1 \wedge w' \neq 0 \wedge \neg u') \end{aligned}$$

Something is wrong. Although $(w=0 \Rightarrow v) \wedge (w=1 \Rightarrow \neg v)$ is a data transformer, it is a rather weak one because when w is neither 0 nor 1 it doesn't constrain v . So the result is that *ask* is transformed into something that's unimplementable.

(c) Replace v with $w: \text{nat}$ according to $(v \Rightarrow w=0) \wedge (\neg v \Rightarrow w=1)$. Is anything wrong?

\S Operation *set* becomes

$$\begin{aligned} & \forall v. (v \Rightarrow w=0) \wedge (\neg v \Rightarrow w=1) \Rightarrow \exists v'. (v' \Rightarrow w'=0) \wedge (\neg v' \Rightarrow w'=1) \wedge (v := \top) \\ & = w: 0,1 \Rightarrow (w := 0) \end{aligned}$$

Operation *flip* becomes

$$\begin{aligned} & \forall v. (v \Rightarrow w=0) \wedge (\neg v \Rightarrow w=1) \Rightarrow \exists v'. (v' \Rightarrow w'=0) \wedge (\neg v' \Rightarrow w'=1) \wedge (v := \neg v) \\ & = w: 0,1 \Rightarrow (w := 1-w) \end{aligned}$$

Operation *ask* becomes

$$\begin{aligned} & \forall v. (v \Rightarrow w=0) \wedge (\neg v \Rightarrow w=1) \Rightarrow \exists v'. (v' \Rightarrow w'=0) \wedge (\neg v' \Rightarrow w'=1) \wedge (u := v) \\ & = w: 0,1 \Rightarrow (u := w=0) \end{aligned}$$

Something is wrong. We have been transforming with something that isn't a transformer; it's too strong.

$$\begin{aligned} & \forall w. \exists v. (v \Rightarrow w=0) \wedge (\neg v \Rightarrow w=1) \\ & = \forall w. w=0 \vee w=1 \\ & = \perp \end{aligned}$$

The last line isn't a theorem, so neither is the first. Nothing constrains the implementation to start in a state where $w=0 \vee w=1$. If it starts with $w=2$, then *set* might not set w to 0, after which *ask* will give the wrong answer.