

CSC C63 Final Exam

April 10, 2015

NAME:

Calculators are not permitted (nor would they be useful).

This is a closed book exam.

#1	/24	#5	/12
#2	/16	#6	/12
#3	/24	#7	/15
#4	/35		
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			/138

You can assume that the following problems are NP-complete.

### CNF-SAT

**Input:** A CNF boolean formula,  $F$ .

**Question:** Is  $F$  satisfiable?

### 3-SAT

**Input:** A CNF boolean formula,  $F$ , in which every clause has exactly 3 literals.

**Question:** Is  $F$  satisfiable?

### IND-SET

**Input:** A graph  $G$ , and an integer  $K$ .

**Question:** Does  $G$  have an independent set of size  $K$ ?

### CLIQUE

**Input:** A graph  $G$ , and an integer  $K$ .

**Question:** Does  $G$  have a clique of size  $K$ ?

### VERTEX-COVER

**Input:** A graph  $G$ , and an integer  $K$ .

**Question:** Does  $G$  have a vertex cover of size  $K$ ?

### 3-COLOUR

**Input:** A graph  $G$ .

**Question:** Does  $G$  have a proper 3-colouring?

### SUBSET-SUM

**Input:** A list of positive integers  $x_1, \dots, x_n$ , and a target  $T$ .

**Question:** Is there a subset of the integers which sums to  $T$ ?

### HAM-PATH

**Input:** A graph  $G$ , with two specified vertices  $u, v$ .

**Question:** Does  $G$  have a Hamilton path from  $u$  to  $v$ ?

### HAM-CYCLE

**Input:** A graph  $G$ .

**Question:** Does  $G$  have a Hamilton cycle?

### KNAPSACK

**Input:** A list of weights  $w_1, \dots, w_n$  and values  $v_1, \dots, v_n$  along with a capacity  $W$  and a target  $T$ . All numbers are positive integers.

**Question:** Is there a subset of indices  $I \subset \{1, \dots, n\}$  such that  $\sum_{i \in I} w_i \leq W$  and  $\sum_{i \in I} v_i \geq T$ ?

1. (24 pts) For each of the following statements, say that it is one of:

- True.
- False.
- No one knows, but most experts think it is True.
- No one knows, but most experts think it is False.

Do not explain your answer.

(a)  $P = NP$

(b)  $PSPACE = NPSPACE$

(c)  $P = EXPTIME$

(d)  $SUBSET-SUM \in co-NP$

(e)  $FACTOR \in P$

(f)  $PATH \in L$

(g) The Post Correspondence Problem is in  $EXPTIME$ .

(h) There is an algorithm to determine whether an algorithm will sort a list of numbers.

2. (16 pts) Only short answers are required here.

(a) (4 pts) State Church's Thesis

(b) (4 pts) Describe the tapes used by a Turing Machine that provides a logspace reduction; i.e. that demonstrates  $A \leq_L B$ .

(c) (4 pts) Give an example of a problem that you know is impossible to solve with an algorithm.

(d) (4 pts)

The following algorithm tests Goldbach's Conjecture for specific numbers. It takes an even number  $2x$  and checks whether it is the sum of two primes.

Input  $x$  (a positive integer)

For  $p = 1$  to  $x$

    Check whether  $p$  is a prime using the polytime prime-testing algorithm

    Check whether  $2x - p$  is a prime using the polytime prime-testing algorithm

    If  $p$  and  $2x - p$  are both prime then Accept

Reject

Is this a polytime algorithm? Explain.

3. (24 pts) Consider the set

$$A = \{ \langle P \rangle, k \mid \text{there is no } x \geq 2k \text{ such that } P(x) \text{ halts and returns } x + k + 1. \}$$

$P$  denotes a Turing Machine and  $k$  is a **positive** integer.

(a) (15 pts) Is  $A$  decidable? Prove your answer.

(b) (3 pts) State  $\bar{A}$  in the form  $\bar{A} = \{\dots\}$ .

(c) (3pts) Is  $\bar{A}$  decidable? Prove your answer.

(d) (3pts) Either  $A$  or  $\bar{A}$  is recognizable. Which one? You do not have to prove your answer.

4. (35 pts) Each of the following problems is in NP; you do not have to prove this (until later in the exam).

For each problem, either prove that it is NP-complete or prove that it is in P.

- (a) (10 pts)

**HUGE-IND-SET**

**Input:** A graph  $G$ .

**Question:** Does  $G$  have an independent set of size  $n - 1$ ?

- (b) (10 pts)

**EVEN-SUBSET-SUM**

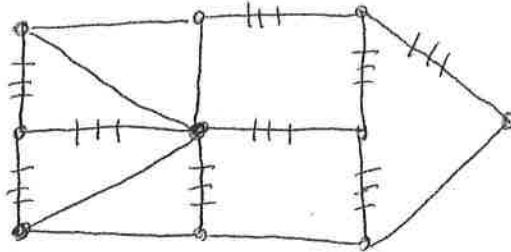
**Input:** A list of positive integers  $x_1, \dots, x_n$ , and a target  $T$  which is an **even number**.

**Question:** Is there a subset of the integers which sums to  $T$ ?



(c) (15 pts)

A  $(1,3)$ -tree is a tree in which every vertex has degree either 1 or 3. A tree  $T$  is a spanning  $(1,3)$ -tree of a graph  $G$  if (i)  $T$  is a  $(1,3)$ -tree ; (ii)  $T$  is a subgraph of  $G$ ; and (iii)  $T$  contains every vertex of  $G$ . For example, the marked edges indicate a spanning  $(1,3)$ -tree in the following graph:



### SPANNING $(1,3)$ -TREE

**Input:** A graph  $G$ .

**Question:** Does  $G$  have a spanning  $(1,3)$ -tree?

**Hint:** This one is NP-complete. Consider a  $(1,3)$ -tree where every degree 3 vertex has at least one degree 1 neighbour.

5. (12 pts) Choose any two problems from the previous question that you think are NP-complete and prove that they are both in NP. For one problem, do this using a verification algorithm. For the other problem, do this using a nondeterministic Turing Machine.

6. (12 pts) A CNF formula  $F$  is *almost satisfiable* if (i) it is not satisfiable and (ii) every clause has the property that if it is removed then the resulting formula is satisfiable. For example, the following formula is almost satisfiable:

$$(x_1 \vee x_2) \wedge (\overline{x_1} \vee x_2) \wedge (\overline{x_2})$$

If we remove the first clause, then  $x_1 = F, x_2 = F$  is a solution; If we remove the second clause, then  $x_1 = T, x_2 = F$  is a solution; if we remove the third clause, then  $x_1 = T, x_2 = T$  is a solution.

#### **ALMOST-SAT**

**Input:** A CNF formula  $F$ .

**Question:** Is  $F$  almost satisfiable?

Show that ALMOST-SAT is in PSPACE.

7. (15 pts)

- (a) (10pts)  $A$  is a recognizable language.  $A^{++}$  is the set of strings formed by concatenating three strings from  $A$ ; i.e.

$$A^{++} = \{a_1a_2a_3 : a_1, a_2, a_3 \in A\}$$

Prove that  $A^{++}$  is recognizable.

(b) (5pts) Is the following statement true? Prove your answer.

If  $A$  is recognizable and  $A \cup B$  is recognizable then  $B$  is recognizable.