

UNIVERSITY OF TORONTO
Faculty of Arts and Science
DECEMBER 2017 EXAMINATIONS

CSC 438H1F/2404H1F

Duration - 3 hours

No Aids Allowed

There are 9 questions worth a total of 100 marks.

Answer all questions on the question paper, using backs of pages for scratch work.

Check that your exam book has 9 pages (including this cover page).

PLEASE COMPLETE THIS SECTION:

Name _____
(Please underline your family name.)

Student Number _____

FOR USE IN MARKING:

1. _____/12

2. _____/12

3. _____/10

4. _____/10

5. _____/10

6. _____/8

7. _____/8

8. _____/12

9. _____/18

Total: _____/100

1. Let \mathcal{L} be a predicate calculus language and let \mathcal{M} be an \mathcal{L} -structure. Let

$$\Sigma = \text{Th}(\mathcal{M}) = \{A \mid A \text{ is an } \mathcal{L}\text{-sentence and } \mathcal{M} \models A\}$$

[7] a) Prove that Σ is a theory (i.e. prove that Σ is closed under logical consequence).

[5] b) Prove that Σ is a complete theory.

[12] 2. The following are the first two Peano Axioms:

P1: $\forall x (sx \neq 0)$

P2: $\forall x \forall y (sx = sy \rightarrow x = y)$

Is it true that $P1, P2 \models \forall x (x = 0 \vee \exists y (x = sy))$?

If true, give a suitable LK proof justifying this (see the next question for the equality axioms) If false, justify by giving a suitable structure.

[10] 3. Give an LK proof of the sequent

$$\forall x(x + 0 = x) \rightarrow \forall x\forall y(x + (y + 0) = x + y)$$

You do not need to put in weakenings or exchanges.

Here are the LK equality axioms:

EL1: $\rightarrow t = t$

EL2: $t = u \rightarrow u = t$

EL3: $t = u, u = v \rightarrow t = v$

EL4: $t_1 = u_1, \dots, t_n = u_n \rightarrow ft_1\dots t_n = fu_1\dots u_n$, for each f in \mathcal{L} , where f is an n -ary function symbol.

EL5: $t_1 = u_1, \dots, t_n = u_n, Pt_1\dots t_n \rightarrow Pu_1\dots u_n$, for each P in \mathcal{L} , where P is an n -ary predicate symbol.

[10]

4. Recall that **TA** (True Arithmetic) is the set of all sentences A in the language $\mathcal{L}_A = [0, s, +, \cdot, =]$ of arithmetic such that A is true in the standard model \mathbb{N} . Suppose that $A(x)$ is a formula of \mathcal{L}_A whose only free variable is x , such that $A(s^n 0)$ is in **TA** for arbitrarily large $n \in \mathbb{N}$. Show that the infinite set of sentences

$$\mathbf{TA} \cup \{A(c), c \neq 0, c \neq s0, c \neq ss0, \dots\}$$

is satisfiable, where c is a new constant.

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- [10] 5. Suppose that $A = \text{range}(f)$ for some computable unary function f . Give a primitive recursive relation $R(x, y)$ such that

$$A = \{x \mid \exists y R(x, y)\}$$

- [8] 6. Let f_1, f_2, f_3, \dots be a list of all total computable functions $f : \mathbb{N} \rightarrow \mathbb{N}$. Define $F(x, y) = f_x(y)$. Prove that F is not computable.

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- [8] 7. Recall that there is a theorem in the Notes that states that every Δ_0 sentence in **TA** is in **RA**. Use this to prove that every $\exists\Delta_0$ sentence in **TA** is in **RA**.

8. Suppose that $A(x)$ is an $\exists\Delta_0$ formula which represents the r.e. set K in \mathbf{PA} .

[2] (a) State what it means for $A(x)$ to represent K in \mathbf{PA} .

[10] (b) Show that there is a consistent extension of Σ of \mathbf{PA} such that $A(x)$ does *not* represent K in Σ .
(Hint: Form Σ by adding a false axiom to \mathbf{PA} .)

[18] 9. We say that a function $f : \mathbb{N} \rightarrow \mathbb{N}$ is *non-decreasing* if $f(x) \leq f(x + 1)$ for all $x \in \mathbb{N}$. Let

$$A = \{x \mid \{x\}_1 \text{ is nondecreasing}\}$$

Is A r.e.? Is A^c r.e.? Justify your answer. (Do not use Rice's Theorem).