

# Skolemization

Conversion of sentences FOL to CNF requires skolemization.

Skolemization: remove existential quantifiers by introducing new function symbols.

How: For each existentially quantified variable introduce a n-place function where n is the number of previously appearing universal quantifiers.

Special case: introducing constants (trivial functions: no previous universal quantifier).

## Skolemization - Example 1

- ▶ Every philosopher writes at least one book.  
 $\forall x[Philo(x) \rightarrow \exists y[Book(y) \wedge Write(x, y)]]$

## Skolemization - Example 1

- ▶ Every philosopher writes at least one book.  
 $\forall x[Philo(x) \rightarrow \exists y[Book(y) \wedge Write(x, y)]]$
- ▶ Eliminate Implication:  
 $\forall x[\neg Philo(x) \vee \exists y[Book(y) \wedge Write(x, y)]]$

## Skolemization - Example 1

- ▶ Every philosopher writes at least one book.  
 $\forall x[Philo(x) \rightarrow \exists y[Book(y) \wedge Write(x, y)]]$
- ▶ Eliminate Implication:  
 $\forall x[\neg Philo(x) \vee \exists y[Book(y) \wedge Write(x, y)]]$
- ▶ Skolemize: substitute  $y$  by  $g(x)$   
 $\forall x[\neg Philo(x) \vee [Book(g(x)) \wedge Write(x, g(x))]]$

## Skolemization - Example 2

- All students of a philosopher read one of their teacher's books.

$$\forall x \forall y [ \text{Philo}(x) \wedge \text{StudentOf}(y, x) \rightarrow \exists z [ \text{Book}(z) \wedge \text{Write}(x, z) \wedge \text{Read}(y, z) ] ]$$

## Skolemization - Example 2

- ▶ All students of a philosopher read one of their teacher's books.

$$\forall x \forall y [Philo(x) \wedge StudentOf(y, x) \rightarrow \exists z [Book(z) \wedge Write(x, z) \wedge Read(y, z)]]$$

- ▶ Eliminate Implication:

$$\forall x \forall y [\neg Philo(x) \vee \neg StudentOf(y, x) \vee \exists z [Book(z) \wedge Write(x, z) \wedge Read(y, z)]]$$

## Skolemization - Example 2

- ▶ All students of a philosopher read one of their teacher's books.

$$\forall x \forall y [\text{Philo}(x) \wedge \text{StudentOf}(y, x) \rightarrow \exists z [\text{Book}(z) \wedge \text{Write}(x, z) \wedge \text{Read}(y, z)]]$$

- ▶ Eliminate Implication:

$$\forall x \forall y [\neg \text{Philo}(x) \vee \neg \text{StudentOf}(y, x) \vee \exists z [\text{Book}(z) \wedge \text{Write}(x, z) \wedge \text{Read}(y, z)]]$$

- ▶ Skolemize: substitute  $z$  by  $h(x, y)$

$$\forall x \forall y [\neg \text{Philo}(x) \vee \neg \text{StudentOf}(y, x) \vee [\text{Book}(h(x, y)) \wedge \text{Write}(x, h(x, y)) \wedge \text{Read}(y, h(x, y))]]$$

## Skolemization - Example 3

- ▶ There exists a philosopher with students.  
 $\exists x \exists y [Philo(x) \wedge StudentOf(y, x)]$

## Skolemization - Example 3

- ▶ There exists a philosopher with students.  
 $\exists x \exists y [Philo(x) \wedge StudentOf(y, x)]$
- ▶ Skolemize: substitute  $x$  by  $a$  and  $y$  by  $b$   
 $Philo(a) \wedge StudentOf(b, a)$

## Most General Unifier

Least specialized unification of two clauses.

We can compute the MGU using the disagreement set

$D_k = \{e_1, e_2\}$ : the pair of expressions where two clauses first disagree.

REPEAT UNTIL no more disagreement  $\rightarrow$  found MGU.

IF either  $e_1$  or  $e_2$  is a variable  $V$  and the other is some term (or a variable)  $t$ , then choose  $V = t$  as substitution.

Then substitute to obtain  $S_{k+1}$  and find disagreement set  $D_{k+1}$ .

ELSE unification is not possible.

## MGU - Example 1

Find the MGU of  $p(f(a), g(X))$  and  $p(Y, Y)$ :

- ▶  $S_0 = \{p(f(a), g(X)) ; p(Y, Y)\}$

## MGU - Example 1

Find the MGU of  $p(f(a), g(X))$  and  $p(Y, Y)$ :

- ▶  $S_0 = \{p(f(a), g(X)) ; p(Y, Y)\}$
- ▶  $D_0 = \{f(a), Y\}$
- ▶  $\sigma = \{Y = f(a)\}$

## MGU - Example 1

Find the MGU of  $p(f(a), g(X))$  and  $p(Y, Y)$ :

- ▶  $S_0 = \{p(f(a), g(X)) ; p(Y, Y)\}$
- ▶  $D_0 = \{f(a), Y\}$
- ▶  $\sigma = \{Y = f(a)\}$
- ▶  $S_1 = \{p(f(a), g(X)) ; p(f(a), f(a))\}$

## MGU - Example 1

Find the MGU of  $p(f(a), g(X))$  and  $p(Y, Y)$ :

- ▶  $S_0 = \{p(f(a), g(X)) ; p(Y, Y)\}$
- ▶  $D_0 = \{f(a), Y\}$
- ▶  $\sigma = \{Y = f(a)\}$
- ▶  $S_1 = \{p(f(a), g(X)) ; p(f(a), f(a))\}$
- ▶  $D_1 = \{g(X), f(a)\}$

## MGU - Example 1

Find the MGU of  $p(f(a), g(X))$  and  $p(Y, Y)$ :

- ▶  $S_0 = \{p(f(a), g(X)) ; p(Y, Y)\}$
- ▶  $D_0 = \{f(a), Y\}$
- ▶  $\sigma = \{Y = f(a)\}$
- ▶  $S_1 = \{p(f(a), g(X)) ; p(f(a), f(a))\}$
- ▶  $D_1 = \{g(X), f(a)\}$
- ▶ no unification possible!

## MGU - Example 2

- ▶  $S_0 = \{p(a, X, h(g(Z))) ; p(Z, h(Y), h(Y))\}$

## MGU - Example 2

- ▶  $S_0 = \{p(a, X, h(g(Z))) ; p(Z, h(Y), h(Y))\}$
- ▶  $D_0 = \{a, Z\}$

## MGU - Example 2

- ▶  $S_0 = \{p(a, X, h(g(Z))) ; p(Z, h(Y), h(Y))\}$
- ▶  $D_0 = \{a, Z\}$
- ▶  $\sigma = \{Z = a\}$

## MGU - Example 2

- ▶  $S_0 = \{p(a, X, h(g(Z))) ; p(Z, h(Y), h(Y))\}$
- ▶  $D_0 = \{a, Z\}$
- ▶  $\sigma = \{Z = a\}$
- ▶  $S_1 = \{p(a, X, h(g(a))) ; p(a, h(Y), h(Y))\}$

## MGU - Example 2

- ▶  $S_0 = \{p(a, X, h(g(Z))) ; p(Z, h(Y), h(Y))\}$
- ▶  $D_0 = \{a, Z\}$
- ▶  $\sigma = \{Z = a\}$
- ▶  $S_1 = \{p(a, X, h(g(a))) ; p(a, h(Y), h(Y))\}$
- ▶  $D_1 = \{X, h(Y)\}$

## MGU - Example 2

- ▶  $S_0 = \{p(a, X, h(g(Z))) ; p(Z, h(Y), h(Y))\}$
- ▶  $D_0 = \{a, Z\}$
- ▶  $\sigma = \{Z = a\}$
- ▶  $S_1 = \{p(a, X, h(g(a))) ; p(a, h(Y), h(Y))\}$
- ▶  $D_1 = \{X, h(Y)\}$
- ▶  $\sigma = \{Z = a, X = h(Y)\}$

## MGU - Example 2

- ▶  $S_0 = \{p(a, X, h(g(Z))) ; p(Z, h(Y), h(Y))\}$
- ▶  $D_0 = \{a, Z\}$
- ▶  $\sigma = \{Z = a\}$
- ▶  $S_1 = \{p(a, X, h(g(a))) ; p(a, h(Y), h(Y))\}$
- ▶  $D_1 = \{X, h(Y)\}$
- ▶  $\sigma = \{Z = a, X = h(Y)\}$
- ▶  $S_2 = \{p(a, h(Y), h(g(a))) ; p(a, h(Y), h(Y))\}$

## MGU - Example 2

- ▶  $S_0 = \{p(a, X, h(g(Z))) ; p(Z, h(Y), h(Y))\}$
- ▶  $D_0 = \{a, Z\}$
- ▶  $\sigma = \{Z = a\}$
- ▶  $S_1 = \{p(a, X, h(g(a))) ; p(a, h(Y), h(Y))\}$
- ▶  $D_1 = \{X, h(Y)\}$
- ▶  $\sigma = \{Z = a, X = h(Y)\}$
- ▶  $S_2 = \{p(a, h(Y), h(g(a))) ; p(a, h(Y), h(Y))\}$
- ▶  $D_2 = \{g(a), Y\}$

## MGU - Example 2

- ▶  $S_0 = \{p(a, X, h(g(Z))) ; p(Z, h(Y), h(Y))\}$
- ▶  $D_0 = \{a, Z\}$
- ▶  $\sigma = \{Z = a\}$
- ▶  $S_1 = \{p(a, X, h(g(a))) ; p(a, h(Y), h(Y))\}$
- ▶  $D_1 = \{X, h(Y)\}$
- ▶  $\sigma = \{Z = a, X = h(Y)\}$
- ▶  $S_2 = \{p(a, h(Y), h(g(a))) ; p(a, h(Y), h(Y))\}$
- ▶  $D_2 = \{g(a), Y\}$
- ▶  $\sigma = \{Z = a, X = h(Y), Y = g(a)\}$

## MGU - Example 2

- ▶  $S_0 = \{p(a, X, h(g(Z))) ; p(Z, h(Y), h(Y))\}$
- ▶  $D_0 = \{a, Z\}$
- ▶  $\sigma = \{Z = a\}$
- ▶  $S_1 = \{p(a, X, h(g(a))) ; p(a, h(Y), h(Y))\}$
- ▶  $D_1 = \{X, h(Y)\}$
- ▶  $\sigma = \{Z = a, X = h(Y)\}$
- ▶  $S_2 = \{p(a, h(Y), h(g(a))) ; p(a, h(Y), h(Y))\}$
- ▶  $D_2 = \{g(a), Y\}$
- ▶  $\sigma = \{Z = a, X = h(Y), Y = g(a)\}$
- ▶  $S_3 = \{p(a, h(g(a)), h(g(a))) ; p(a, h(g(a)), h(g(a)))\}$

## MGU - Example 2

- ▶  $S_0 = \{p(a, X, h(g(Z))) ; p(Z, h(Y), h(Y))\}$
- ▶  $D_0 = \{a, Z\}$
- ▶  $\sigma = \{Z = a\}$
- ▶  $S_1 = \{p(a, X, h(g(a))) ; p(a, h(Y), h(Y))\}$
- ▶  $D_1 = \{X, h(Y)\}$
- ▶  $\sigma = \{Z = a, X = h(Y)\}$
- ▶  $S_2 = \{p(a, h(Y), h(g(a))) ; p(a, h(Y), h(Y))\}$
- ▶  $D_2 = \{g(a), Y\}$
- ▶  $\sigma = \{Z = a, X = h(Y), Y = g(a)\}$
- ▶  $S_3 = \{p(a, h(g(a)), h(g(a))) ; p(a, h(g(a)), h(g(a)))\}$
- ▶ **No disagreement**  
 $\Rightarrow \sigma = \{Z = a, X = h(Y), Y = g(a)\}$  is MGU

## MGU - Example 3

- ▶  $S_0 = \{p(X, X) ; p(Y, f(Y))\}$

## MGU - Example 3

- ▶  $S_0 = \{p(X, X) ; p(Y, f(Y))\}$
- ▶  $D_0 = \{X, Y\}$

## MGU - Example 3

- ▶  $S_0 = \{p(X, X) ; p(Y, f(Y))\}$
- ▶  $D_0 = \{X, Y\}$
- ▶  $\sigma = \{X = Y\}$

## MGU - Example 3

- ▶  $S_0 = \{p(X, X) ; p(Y, f(Y))\}$
- ▶  $D_0 = \{X, Y\}$
- ▶  $\sigma = \{X = Y\}$
- ▶  $S_1 = \{p(Y, Y) ; p(Y, f(Y))\}$

## MGU - Example 3

- ▶  $S_0 = \{p(X, X) ; p(Y, f(Y))\}$
- ▶  $D_0 = \{X, Y\}$
- ▶  $\sigma = \{X = Y\}$
- ▶  $S_1 = \{p(Y, Y) ; p(Y, f(Y))\}$
- ▶  $D_1 = \{Y, f(Y)\}$

## MGU - Example 3

- ▶  $S_0 = \{p(X, X) ; p(Y, f(Y))\}$
- ▶  $D_0 = \{X, Y\}$
- ▶  $\sigma = \{X = Y\}$
- ▶  $S_1 = \{p(Y, Y) ; p(Y, f(Y))\}$
- ▶  $D_1 = \{Y, f(Y)\}$
- ▶ no unification possible!

## Full example problem

Given the following sentences, answer the question '**What is connected to the Galbraith building?**' using resolution with answer extraction:

Connected is a binary symmetric relation.

An object  $X$  is part of another object  $Y$  iff everything  $X$  is connected to,  $Y$  is also connected to.

Room GB221 is part of Galbraith building.

Room GB221 is connected to itself.

## Full example problem - Representation in FOL

(a) Represent these sentences in first order logic.

Connected is a binary symmetric relation.

An object  $X$  is part of another object  $Y$  iff everything  $X$  is connected to,  $Y$  is also connected to.

Room GB221 is part of Galbraith building.

Room GB221 is connected to itself.

## Full example problem - Representation in FOL

- ▶ Connected is a symmetric relation.

$$(\forall X, Y) \text{connected}(X, Y) \supset \text{connected}(Y, X)$$

## Full example problem - Representation in FOL

- ▶ An object  $X$  is part of another object  $Y$  iff everything  $X$  is connected to,  $Y$  is also connected to.

$$(\forall X, Y) (\text{part}(X, Y) \equiv ((\forall Z) \text{connected}(Z, X) \supset \text{connected}(Z, Y)))$$

## Full example problem - Representation in FOL

- ▶ Room GB221 is part of Galbraith building.  
 $part(gb221, galbraith)$

## Full example problem - Representation in FOL

- ▶ Room GB221 is connected to itself.  
 $connected(gb221, gb221)$

## Full example problem - CNF Conversion

(b) Convert the formulas to clausal form. Indicate any Skolem functions or constants used.

$(\forall X, Y) \text{connected}(X, Y) \supset \text{connected}(Y, X)$

$(\forall X, Y) \text{part}(X, Y) \equiv (\forall Z) \text{connected}(Z, X) \supset \text{connected}(Z, Y)$

$\text{part}(\text{gb221}, \text{galbraith})$

$\text{connected}(\text{gb221}, \text{gb221})$

## Full example problem - CNF Conversion

- ▶  $(\forall X, Y) \text{connected}(X, Y) \supset \text{connected}(Y, X)$   
[ $\neg \text{connected}(X, Y), \text{connected}(Y, X)$ ]

## Full example problem - CNF Conversion

- ▶  $(\forall X, Y) part(X, Y) \equiv (\forall Z) connected(Z, X) \supset connected(Z, Y)$   
→:  $[-part(X, Y), -connected(Z, X), connected(Z, Y)]$   
←:  $[part(X, Y), connected(f(X, Y), X)]$   
 $[part(X, Y), -connected(g(X, Y), Y)]$

## Full example problem - CNF Conversion

- ▶  $\text{part}(gb221, \text{galbraith})$   
 $[\text{part}(gb221, \text{galbraith})]$

## Full example problem - CNF Conversion

- ▶  $\text{connected}(gb221, gb221)$   
 $[\text{connected}(gb221, gb221)]$

## Full example problem - Goal

(c) Convert the negation of the statement '**What is connected to the Galbraith building?**' to clause form (using an answer literal).

- ▶ FOL:  $(\exists X) \text{connected}(\text{galbraith}, X)$

## Full example problem - Goal

(c) Convert the negation of the statement '**What is connected to the Galbraith building?**' to clause form (using an answer literal).

- ▶ FOL:  $(\exists X) \text{connected}(\text{galbraith}, X)$
- ▶ negate goal!!  $(\neg \exists X) \text{connected}(\text{galbraith}, X)$

## Full example problem - Goal

(c) Convert the negation of the statement '**What is connected to the Galbraith building?**' to clause form (using an answer literal).

- ▶ FOL:  $(\exists X) \text{connected}(\text{galbraith}, X)$
- ▶ negate goal!!  $(\neg \exists X) \text{connected}(\text{galbraith}, X)$
- ▶ CNF with answer literal:  $(\forall X) \neg \text{connected}(\text{galbraith}, X)$   
[ $\neg \text{connected}(\text{galbraith}, X), \text{ans}(X)$ ]

## Full example problem - Resolution

(d) Answer the question using resolution and answer extraction.  
Use the notation developed in class: every new clause must be labeled by the resolution step that was used to generate it. For example, a clause labeled  $R[4c, 1d]x = a, y = f(b)$  means that it was generated by resolving literal c of clause 4 against literal d of clause 1, using the MGU  $x = a, y = f(b)$ .

Our clauses:

1.  $[\neg \text{connected}(R, S), \text{connected}(S, R)]$
2.  $[\neg \text{part}(T, U), \neg \text{connected}(V, T), \text{connected}(V, U)]$
3.  $[\text{part}(W, X), \text{connected}(f(W, X), W)]$
4.  $[\text{part}(Y, Z), \neg \text{connected}(g(Y, Z), Z)]$
5.  $[\text{part}(gb221, galbraith)]$
6.  $[\text{connected}(gb221, gb221)]$
7.  $[\neg \text{connected}(galbraith, A), \text{ans}(A)]$

## Full example problem - Resolution

1.  $[\neg \text{connected}(R, S), \text{connected}(S, R)]$
2.  $[\neg \text{part}(T, U), \neg \text{connected}(V, T), \text{connected}(V, U)]$
3.  $[\text{part}(W, X), \text{connected}(f(W, X), W)]$
4.  $[\text{part}(Y, Z), \neg \text{connected}(g(Y, Z), Z)]$
5.  $[\text{part}(gb221, galbraith)]$
6.  $[\text{connected}(gb221, gb221)]$
7.  $[\neg \text{connected}(galbraith, A), \text{ans}(A)]$

## Full example problem - Resolution

1.  $[\neg \text{connected}(R, S), \text{connected}(S, R)]$
2.  $[\neg \text{part}(T, U), \neg \text{connected}(V, T), \text{connected}(V, U)]$
3.  $[\text{part}(W, X), \text{connected}(f(W, X), W)]$
4.  $[\text{part}(Y, Z), \neg \text{connected}(g(Y, Z), Z)]$
5.  $[\text{part}(gb221, galbraith)]$
6.  $[\text{connected}(gb221, gb221)]$
7.  $[\neg \text{connected}(galbraith, A), \text{ans}(A)]$
8. R[7a, 1b]  $\{S = galbraith, R = U\}$   
 $[\neg \text{connected}(A, galbraith), \text{ans}(A)]$

## Full example problem - Resolution

1.  $[\neg \text{connected}(R, S), \text{connected}(S, R)]$
2.  $[\neg \text{part}(T, U), \neg \text{connected}(V, T), \text{connected}(V, U)]$
3.  $[\text{part}(W, X), \text{connected}(f(W, X), W)]$
4.  $[\text{part}(Y, Z), \neg \text{connected}(g(Y, Z), Z)]$
5.  $[\text{part}(gb221, galbraith)]$
6.  $[\text{connected}(gb221, gb221)]$
7.  $[\neg \text{connected}(galbraith, A), \text{ans}(A)]$
8. R[7a, 1b]  $\{S = galbraith, R = U\}$   
 $[\neg \text{connected}(A, galbraith), \text{ans}(A)]$
9. R[8a, 2c]  $\{V = A, U = galbraith\}$   
 $[\neg \text{part}(T, galbraith), \neg \text{connected}(A, T), \text{ans}(A)]$

## Full example problem - Resolution

1.  $[\neg \text{connected}(R, S), \text{connected}(S, R)]$
2.  $[\neg \text{part}(T, U), \neg \text{connected}(V, T), \text{connected}(V, U)]$
3.  $[\text{part}(W, X), \text{connected}(f(W, X), W)]$
4.  $[\text{part}(Y, Z), \neg \text{connected}(g(Y, Z), Z)]$
5.  $[\text{part}(gb221, galbraith)]$
6.  $[\text{connected}(gb221, gb221)]$
7.  $[\neg \text{connected}(galbraith, A), \text{ans}(A)]$
8. R[7a, 1b]  $\{S = galbraith, R = U\}$   
 $[\neg \text{connected}(A, galbraith), \text{ans}(A)]$
9. R[8a, 2c]  $\{V = A, U = galbraith\}$   
 $[\neg \text{part}(T, galbraith), \neg \text{connected}(A, T), \text{ans}(A)]$
10. R[9a, 5]  $\{T = gb221\}$   
 $[\neg \text{connected}(A, gb221), \text{ans}(A)]$

## Full example problem - Resolution

1.  $[\neg \text{connected}(R, S), \text{connected}(S, R)]$
2.  $[\neg \text{part}(T, U), \neg \text{connected}(V, T), \text{connected}(V, U)]$
3.  $[\text{part}(W, X), \text{connected}(f(W, X), W)]$
4.  $[\text{part}(Y, Z), \neg \text{connected}(g(Y, Z), Z)]$
5.  $[\text{part}(gb221, galbraith)]$
6.  $[\text{connected}(gb221, gb221)]$
7.  $[\neg \text{connected}(galbraith, A), \text{ans}(A)]$
8. R[7a, 1b]  $\{S = galbraith, R = U\}$   
 $[\neg \text{connected}(A, galbraith), \text{ans}(A)]$
9. R[8a, 2c]  $\{V = A, U = galbraith\}$   
 $[\neg \text{part}(T, galbraith), \neg \text{connected}(A, T), \text{ans}(A)]$
10. R[9a, 5]  $\{T = gb221\}$   
 $[\neg \text{connected}(A, gb221), \text{ans}(A)]$
11. R[10a, 6]  $\{A = gb221\}$   
 $[\text{ans}(gb221)]$