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Assignment 1 [SOLUTIONS]

Marked Questions

Q1. [10 Points] Fibonacci Numbers (maximum 2 pages)

1. Let f_n be the n^{th} Fibonacci Number where $f_0 = 0$ and $f_1 = 1$. Prove that $\sum_{k=0}^n f_k = f_{n+2} - 1$ for all $n \in \mathbb{N}$.

Proof. Define predicate $P(n) := \sum_{k=0}^{n} f_k = f_{n+2} - 1$. Prove P(n) is true $\forall n \in \mathbb{N}$ by induction on n.

- (a) Base case. Note that $f_0 = 0$ and $f_1 = 1$ (given) as well as $f_2 = 1$ and $f_3 = 2$ (deduced from the definition $f_{n+2} = f_{n+1} + f_n$. Consider the cases where n = 0 and n = 1. In the former case $\sum_{k=0}^{0} f_k = f_0 = 0 = 1 1 = f_2 1$. In the latter case, $\sum_{k=0}^{1} f_k = f_0 + f_1 = 1 = 2 1 = f_3 1$.
- (b) Inductive step. For a fixed m suppose that P(0), ..., P(m) are all true, we want to show that P(m+1) is true. Note that

$$\sum_{k=0}^{n+1} f_k = \left(\sum_{k=0}^m f_k\right) + f_{m+1}$$

= $f_{m+2} - 1 + f_{m+1}$ (IH)
= $f_{m+3} - 1$ (Definition of f_{m+3})

It follow by the Principle of Mathematical Induction that P(n) is true $\forall n \in \mathbb{N}$.

Total: 5 points. [1 Point] Predicate/tell me what you are inducting on. [2 Points] Base case must have at least two cases (one point per case). [2 Points] Inductive step must use the inductive hypothesis in order to get any points. Minor and major arithmetic errors are (-0.5) and (-1) respectively.

2. Prove that $f_n \leq \left(\frac{7}{4}\right)^{n-1}$ for all $n \in \mathbb{N}$ when $n \geq 1$ from first principles (do not used the equation for f_n shown in class).

Proof. Define $P(n) \coloneqq f_n \leq (7/4)^{n-1}$. Prove P(n) is true $\forall n \in \mathbb{N}, n \geq 1$ by induction on n.

- (a) Base case. Consider the cases where n = 1 and n = 2. In the former, we have $f_1 = 1 \le (7/4)^0 = 1$. In the latter case, we have $f_2 = 1 \le (7/4)^1$.
- (b) Inductive step. Fix some $k \ge 2$ and suppose that P(1), P(2), ..., P(k) are true. Show

that P(k+1) is true. Note that

$$f_{k+1} = f_k + f_{k-1} \qquad \text{(Definition of } f_{k+1}\text{)}$$

$$\leq \left(\frac{7}{4}\right)^{k-1} + \left(\frac{7}{4}\right)^{k-2} \qquad \text{(IH)}$$

$$\leq \left(\frac{7}{4}\right)^{k-2} \left(\frac{7}{4} + 1\right) = \left(\frac{7}{4}\right)^{k-2} \left(\frac{44}{16}\right)$$

$$\leq \left(\frac{7}{4}\right)^{k-2} \left(\frac{49}{16}\right) = \left(\frac{7}{4}\right)^k.$$

It follow by the Principle of Mathematical Induction that P(n) is true for all $n \in \mathbb{N}$, $n \ge 1$. \Box

Total: 5 points. [1 Point] Predicate/tell me what you are inducting on. [1 Point] Base case must have at least two cases (0.5 points per case) as this was pretty much already shown in Lecture 1. [3 Points] Inductive step must use the inductive hypothesis in order to get any points. Minor and major arithmetic errors are (-0.5) and (-1) respectively.

Q2. [10 Points] Marshmallow Game (maximum 2 pages)

Alice and Bob's friend made a large pan of marshmallows which are cut into $m \times n$ squares but still stuck together. Alice and Bob decides to play a game with the marshmallow to pull it apart. On her turn, Alice pulls the marshmallow apart along a vertical cut while, on his turn, Bob pulls the marshmallow apart along a horizontal cut. The player pulling the marshmallow then choose one of the pieces to be the *active* and the game continues on the active piece. The first player unable to make a move loses. Alice goes first. Sees a sample game in Figure 1.

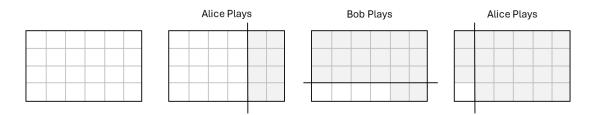


Figure 1: Initially the full grid is active. Alice pulls off two columns of marshmallows. The remaining 4×4 grid is active (the gray portion is inactive). Bob pulls off three rows leaving a 1×4 active grid. Alice pulls off three columns leaving one square. Since Bob cannot move, Alice wins!

1. If m = 1, who is the winner? What about when n = 1?

When $m \ge 1$ and n = 1, Bob is the winner. In particular if m = n = 1, then on the first turn Alice has no valid moves and loses (Bob wins). Generally when n = 1, Alice has no valid moves and loses (Bob wins). Conversely, if m = 1 and n > 1, then Alice is the winner. On her first turn, Alice can pull apart the marshmallow into a single square and another rectangle which has dimension $1 \times (n - 1)$. If she designates the single square to be active, then on Bob's turn he will have no valid moves and loses (Alice wins).

Total: 3 points. [1 Point] Consider case m = 1 and n = 1. [1 Point] Consider case m > 1 and n = 1. [1 Point] Consider case m = 1 and n > 1.

2. For $m, n \in \mathbb{N}$ with $m, n \ge 1$ who wins? Prove your claim.

Claim 1. When $m \ge n$, Bob win. Otherwise Alice wins.

Proof. We prove the following predicate P(m, n) for $m, n \in \mathbb{N}$ with $m, n \geq 1$.

Case 1. If Alice plays first and $m \ge n$ then Bob wins, otherwise Alice wins.

Case 2. If Bob plays first and $n \ge m$ then Alice wins, otherwise Bob wins.

If we show P(m,n) is true $\forall m, n \in \mathbb{N}$ with $m, n \geq 1$, then the claim is true; The first case of the predicate is *exactly* our game (Alice plays first). Apply induction on the pair (m, n).

(a) Base case:

Alice first. This is *exactly* what was considered in the first part of Q2 (above).

Bob first. This is very similar to the first part of Q2 (above). For completeness: when $n \ge 1$ and m = 1, Alice wins. In particular if m = n = 1, then on the first turn Bob has no valid moves and loses. Generally, when m = 1, Bob has no valid moves and loses.

Conversely, if n = 1 and m > 1, then Bob wins. On his first turn, Bob can pull apart the marshmallow into a single square and another rectangle with dimension $(m-1) \times 1$. If he designates the single square to be active, then on Alice's turn she will have no valid moves and loses.

(b) Inductive step: Fix some k > 1 and $\ell > 1$. Suppose $\forall k' : 1 \le k' < k$ and $\forall \ell' :\le \ell' < \ell$, $P(k', \ell), P(k, \ell')$, and $P(k', \ell')$ are true. We want to show that $P(k, \ell)$ is true.

Alice first. Either $\ell > k$ or $\ell \leq k$.

When $\ell > k$, Alice can pull away $\ell - k$ columns to obtain a $k \times k$ square which she will designate as active. Then, Bob will be playing first on a $k \times k$ square. By IH applied to P(k, k), when Bob plays first Alice wins.

When $\ell \leq k$, regardless of how Alice pulls apart the marshmallow, the two resulting rectangles will have dimension $k \times \ell_1$ and $k \times \ell_2$ where $\ell_1 + \ell_2 = \ell$. Note $\ell_1, \ell_2 < \ell \leq k$. By IH applied to $P(k, \ell_1)$ and $P(k, \ell_2)$, when Bob plays first on either of these rectangles, he will win.

Bob first. Very similar to previous case. The following is for completeness only. Either $k > \ell$ or $k \le \ell$.

When $k > \ell$, Bob can pull away $k - \ell$ rows to obtain a $\ell \times \ell$ square which he will designate as active. Then, Alice will be playing first on a $\ell \times \ell$ square. By IH applied to $P(\ell, \ell)$, when Alice plays first Bob wins.

When $k \leq \ell$, regardless of how Bob pulls apart the marshmallow, the two resulting rectangles will have dimension $k_1 \times \ell$ and $k_2 \times \ell$ where $k_1 + k_2 = k$. Note $k_1, k_2 < k \leq \ell$. By IH applied to $P(k_1, \ell)$ and $P(k_2, \ell)$, when Alice plays first on either of these rectangles, she will win.

Thus, by complete induction, P(m, n) is true for all $m, n \in \mathbb{N}$ where $m, n \ge 1$.

Total: 7 points. [3 Points] Predicate. [1 Points] Base case. [3 Points] Inductive step must use IH to get points. Minor and major logical errors get a reduction of -0.5 and -1 respectively. If the predicate has significant errors then this part can get at most two points for the predicate and nothing else. Be careful when marking this, there may be multiple valid solutions e.g. with different number of variables in the predicate, using Well-ordering principle (though induction was encouraged), etc.