**1.** Use the Pumping Lemma to prove that the following language is not regular:

$$L = \{0^i : i \text{ is a prime}\}\$$

Solution

Assume for a contradiction that *L* is regular. Then *L* must satisfy the Pumping Lemma. That is, there exists  $p \in \mathbb{N}$ , p > 0 such that for every  $w \in L$  with  $|w| \ge p$ , conditions (1) to (4) the in the Pumping Lemma are satisfied.

Let  $w = 0^m$  where  $m \ge p + 2$  is prime. Then  $w \in L$  and  $|w| \ge p$ . Assume w = xyz. We show that for all x, y, z, as least one of the conditions (2)–(4) of the Pumping Lemma is not satisfied.

Assume that Conditions (2) and (3) are satisfied; that is,  $|y| \ge 1$  and  $|xy| \le p$ .

Note that since  $w = xyz = 0^m$  (where  $m \ge p + 2$ ) and  $|xy| \le p$ , |z| must be strictly greater than 1, and therefore |xz| > 1.

Now choose i = |xz|. Then,  $|xy^i z| = |xz| + |y||xz| = (1 + |y|)|xz|$ .

Since (1 + |y|) and |xz| are each greater than 1, the product must be a composite number. That is,  $|xy^iz|$  is a composite number and so  $xy^iz$  is not a member of *L*; hence Condition (4) of the Pumping Lemma fails.

This is, however, a contradiction! So we can conclude that the original assumption (that L is regular) is false, and L is non-regular.