

Function Fine Points

partial:

sometimes no result

total:

always at least one result

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total:

always at least one result

deterministic:

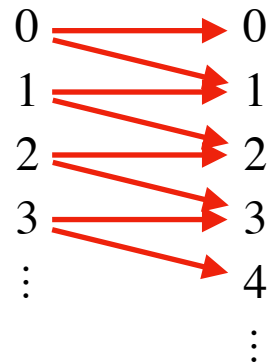
always at most one result

nondeterministic:

sometimes more than one result

Function Fine Points

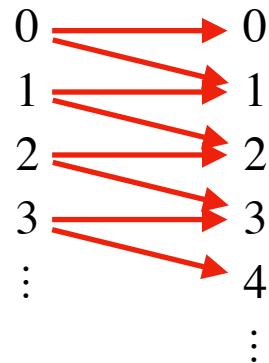
partial:	sometimes no result
total:	always at least one result
deterministic:	always at most one result
nondeterministic:	sometimes more than one result
$\langle n: \text{nat} \cdot n, n+1 \rangle$	total and nondeterministic



Function Fine Points

partial:	sometimes no result
total:	always at least one result
deterministic:	always at most one result
nondeterministic:	sometimes more than one result

$\langle n: \text{nat} \cdot n, n+1 \rangle$ total and nondeterministic



$\langle n: \text{nat} \cdot n, n+1 \rangle 3 = 3, 4$

distribution

$$(f, g) x = f x, g x$$

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$$f(x, y) = f x, f y$$

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$$double = \langle n: nat \cdot n+n \rangle$$

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$$\text{double} = \langle n: \text{nat} \cdot n+n \rangle$$

$$\text{double } 2 = 4$$

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$$\text{double } 2 = 4$$

$$\text{double } (2, 3) = \text{double } 2, \text{double } 3 = 4, 6$$

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$$\text{tiny} = \langle S: \text{nat} \cdot S < 3 \rangle$$

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$$\text{tiny } \{\text{null}\} = \top$$

$$\text{tiny } \{0, 1, 2, 3\} = \perp$$

$$\text{tiny } \text{null} = \text{null}$$

function inclusion

$$f: g = \Box f :: \Box g \wedge \forall x: \Box g \cdot f x: g x$$

function inclusion

$$f: g = \Box f :: \Box g \wedge \forall x: \Box g \cdot f x: g x$$

$$f: A \rightarrow B$$

function inclusion

$$f: g = \Box f :: \Box g \wedge \forall x: \Box g \cdot f x: g x$$

$$A \rightarrow B = \langle a: A \cdot B \rangle$$

$$f: A \rightarrow B$$

function inclusion

$$f: g = \Box f :: \Box g \wedge \forall x: \Box g \cdot f x: g x$$

$$A \rightarrow B = \langle a: A \cdot B \rangle$$

$A \rightarrow B$ is a function whose domain is A and whose result is B .

$$f: A \rightarrow B$$

function inclusion

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$$f: A \rightarrow B = \Box f:: A \wedge \forall a: A \cdot f a: B$$

$A \rightarrow B$ is all those functions whose domain is at least A and whose result is at most B .

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$A \rightarrow B$ is antimonotonic in A and monotonic in B .

function inclusion

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$$\text{*suc*: nat} \rightarrow \text{*nat*}$$

function inclusion

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$$suc: nat \rightarrow nat$$

function inclusion

$$= \Box suc:: nat \wedge \forall n: nat \cdot suc n: nat$$

function inclusion

$$f: g = \Box f:: \Box g \wedge \forall x: \Box g \cdot f x: g x$$

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$$suc: nat \rightarrow nat$$

function inclusion

$$= \Box suc:: nat \wedge \forall n: nat \cdot suc n: nat$$

definition of *suc*

$$= nat:: nat \wedge \forall n: nat \cdot n+1: nat$$

function inclusion

$$f: g = \Box f:: \Box g \wedge \forall x: \Box g \cdot f x: g x$$

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$$suc: nat \rightarrow nat$$

function inclusion

$$= \Box suc:: nat \wedge \forall n: nat \cdot suc n: nat$$

definition of suc

$$= nat:: nat \wedge \forall n: nat \cdot n+1: nat$$

reflexivity and definition of nat

$$= \top$$

function inclusion

$$f: g = \Box f :: \Box g \wedge \forall x: \Box g \cdot f x: g x$$

function inclusion

$$f: g = \Box f :: \Box g \wedge \forall x: \Box g \cdot f x: g x$$

suc: nat → nat

even: int → bin

avg: rat → rat → rat

function inclusion

$$f: g = \Box f :: \Box g \wedge \forall x: \Box g \cdot f x: g x$$

suc: nat → nat

even: int → bin

avg: rat → rat → rat

$$A: B \wedge f: B \rightarrow C \wedge C: D \Rightarrow f: A \rightarrow D$$

function inclusion

$$f: g = \Box f :: \Box g \wedge \forall x: \Box g \cdot f x: g x$$

suc: nat → nat

even: int → bin

avg: rat → rat → rat

$$A: B \wedge f: B \rightarrow C \wedge C: D \Rightarrow f: A \rightarrow D$$

$$(0,..10): nat \wedge suc: nat \rightarrow nat \wedge nat: int \Rightarrow suc: (0,..10) \rightarrow int$$

function inclusion

$$f: g = \Box f :: \Box g \wedge \forall x: \Box g \cdot f x: g x$$

suc: nat → nat

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$$A: B \wedge f: B \rightarrow C \wedge C: D \Rightarrow f: A \rightarrow D$$

$$(0,..10): nat \wedge suc: nat \rightarrow nat \wedge nat: int \Rightarrow suc: (0,..10) \rightarrow int$$

$$\langle f: (0,..10) \rightarrow int \cdot \forall n: 0,..10 \cdot even (f n) \rangle$$

function inclusion

$$f: g = \Box f :: \Box g \wedge \forall x: \Box g \cdot f x: g x$$

suc: nat → nat

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$$(0,..10): nat \wedge suc: nat \rightarrow nat \wedge nat: int \Rightarrow suc: (0,..10) \rightarrow int$$

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$$\langle f: (0,..10) \rightarrow int \cdot \underline{\forall n: 0,..10} \cdot \underline{even (f n)} \rangle$$

function inclusion

$$f: g = \Box f :: \Box g \wedge \forall x: \Box g \cdot f x: g x$$

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$$A: B \wedge f: B \rightarrow C \wedge C: D \Rightarrow f: A \rightarrow D$$

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$$\langle f: (0,..10) \rightarrow int \cdot \forall n: 0,..10 \cdot \underline{even (f n)} \rangle$$

function inclusion

$$f: g = \Box f :: \Box g \wedge \forall x: \Box g \cdot f x: g x$$

suc: nat → nat

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avg: rat → rat → rat

$$A: B \wedge f: B \rightarrow C \wedge C: D \Rightarrow f: A \rightarrow D$$

$$(0, \dots, 10): \text{nat} \wedge \text{suc}: \text{nat} \rightarrow \text{nat} \wedge \text{nat}: \text{int} \Rightarrow \text{suc}: (0, \dots, 10) \rightarrow \text{int}$$

$$\langle f: (0, \dots, 10) \rightarrow \text{int} \cdot \forall n: 0, \dots, 10 \cdot \text{even} (f n) \rangle \text{suc}$$

function inclusion

$$f: g = \Box f :: \Box g \wedge \forall x: \Box g \cdot f x: g x$$

$$suc: nat \rightarrow nat$$

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$$(0,..10): nat \wedge suc: nat \rightarrow nat \wedge nat: int \Rightarrow suc: (0,..10) \rightarrow int$$

$$\langle f: (0,..10) \rightarrow int \cdot \forall n: 0,..10 \cdot even (f n) \rangle suc$$

$$= \forall n: 0,..10 \cdot even (suc n)$$

function inclusion

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$$(0,..10): nat \wedge suc: nat \rightarrow nat \wedge nat: int \Rightarrow suc: (0,..10) \rightarrow int$$

$$\langle f: (0,..10) \rightarrow int \cdot \forall n: 0,..10 \cdot even (f n) \rangle suc$$

$$= \forall n: 0,..10 \cdot even (suc n)$$

$$= \perp$$

function equality

$$f = g \quad = \quad \Box f = \Box g \quad \wedge \quad \forall x: \Box f \cdot f x = g x$$

function composition

If $\neg f: \square g$ then

$$\square(g f) = \S x: \square f f x: \square g$$

$$(g f) x = g (f x)$$

function composition

If $\neg f: \square g$ then

$$\square(g f) = \S x: \square f \cdot f x: \square g$$

$$(g f) x = g (f x) \quad \leftarrow$$

function composition

If $\neg f: \square g$ then

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$$(g f) x = g (f x)$$

function composition

If $\neg f: \square g$ then

$$\square(g f) = \S x: \square f \cdot f x: \square g$$

$$(g f) x = g (f x)$$

$$\square(\text{even suc})$$

$$= \S x: \square \text{suc} \cdot \text{suc } x: \square \text{even}$$

$$= \S x: \text{nat} \cdot x+1: \text{int}$$

$$= \text{nat}$$

function composition

If $\neg f: \square g$ then

$$\square(g f) = \S x: \square f \cdot f x: \square g$$

$$(g f) x = g (f x)$$

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$$= \S x: \square \text{suc} \cdot \text{suc } x: \square \text{even}$$

$$= \S x: \text{nat} \cdot x+1: \text{int}$$

$$= \text{nat}$$

$$(\text{even suc}) 3$$

function composition

If $\neg f: \square g$ then

$$\square(g f) = \S x: \square f \cdot f x: \square g$$

$$(g f) x = g (f x)$$

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$$= \S x: \square \text{suc} \cdot \text{suc } x: \square \text{even}$$

$$= \S x: \text{nat} \cdot x+1: \text{int}$$

$$= \text{nat}$$

$$(\text{even suc}) 3 = \text{even} (\text{suc } 3)$$

function composition

If $\neg f: \square g$ then

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$$= \S x: \text{nat} \cdot x+1: \text{int}$$

$$= \text{nat}$$

$$(\text{even suc}) 3 = \text{even} (\text{suc } 3) = \text{even } 4$$

function composition

If $\neg f: \square g$ then

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$$(g f) x = g (f x)$$

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$$= \S x: \square \text{suc} \cdot \text{suc } x: \square \text{even}$$

$$= \S x: \text{nat} \cdot x+1: \text{int}$$

$$= \text{nat}$$

$$(\text{even suc}) 3 = \text{even} (\text{suc } 3) = \text{even } 4 = \top$$

function composition

If $\neg f: \square g$ then

$$\square(g f) = \S x: \square f \cdot f x: \square g$$

$$(g f) x = g (f x)$$

$$\square(\text{even suc})$$

$$= \S x: \square \text{suc} \cdot \text{suc } x: \square \text{even}$$

$$= \S x: \text{nat} \cdot x+1: \text{int}$$

$$= \text{nat}$$

$$(\text{even suc}) 3 = \text{even} (\text{suc } 3) = \text{even } 4 = \top$$

$$(-\text{suc}) 3 = -(\text{suc } 3) = -4$$

function composition

If $\neg f: \square g$ then

$$\square(g f) = \S x: \square f \cdot f x: \square g$$

$$(g f) x = g (f x)$$

$$\square(\text{even suc})$$

$$= \S x: \square \text{suc} \cdot \text{suc } x: \square \text{even}$$

$$= \S x: \text{nat} \cdot x+1: \text{int}$$

$$= \text{nat}$$

$$(\text{even suc}) 3 = \text{even} (\text{suc } 3) = \text{even } 4 = \top$$

$$(-\text{suc}) 3 = -(\text{suc } 3) = -4$$

$$(\neg \text{even}) 3 = \neg(\text{even } 3) = \neg \perp = \top$$

function composition

function composition

Suppose $x, y: int$

$f, g: int \rightarrow int$

$h: int \rightarrow int \rightarrow int$

function composition

Suppose $x, y: int$

$f, g: int \rightarrow int$

$h: int \rightarrow int \rightarrow int$

Then

$h f x g y$

function composition

Suppose $x, y: int$

$f, g: int \rightarrow int$

$h: int \rightarrow int \rightarrow int$

Then

$$\begin{aligned} & h f x g y \\ = & ((h f) x) g y \end{aligned}$$

function composition

Suppose $x, y: int$

$f, g: int \rightarrow int$

$h: int \rightarrow int \rightarrow int$

Then

$$\begin{aligned} & h f x g y \\ = & ((h f) x) g y \\ = & (h (f x)) g y \end{aligned}$$

function composition

Suppose $x, y: int$

$f, g: int \rightarrow int$

$h: int \rightarrow int \rightarrow int$

Then

$$\begin{aligned} & h f x g y \\ = & ((h f) x) g y \\ = & (h (f x)) g y \\ = & (h (f x)) (g y) \end{aligned}$$

function composition

Suppose $x, y: int$

$f, g: int \rightarrow int$

$h: int \rightarrow int \rightarrow int$

Then

$$\begin{aligned} & h f x g y \\ = & (((h f) x) g) y \\ = & ((h (f x)) g) y \\ = & (h (f x)) (g y) \\ = & h (f x) (g y) \end{aligned}$$

list as function

If $m: \square L$ then

$$\langle n: \square L \cdot L n \rangle$$

list as function

If $m: \square L$ then

$$\langle n: \square L \cdot L n \rangle m$$

list as function

If $m: \square L$ then

$$\langle n: \square L \cdot L n \rangle m = L m$$

list as function

If $m: \square L$ then

$$\langle \underline{n: \square L \cdot L n} \rangle m = \underline{L m}$$

function \approx list

application \approx indexing

list as function

If $m: \square L$ then

$$\langle n: \square L \cdot L n \rangle m = L m$$

function \approx list

application \approx indexing

function composition \approx list composition

list as function

If $m: \square L$ then

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function \approx list

application \approx indexing

function composition \approx list composition

– [3; 5; 2]

suc [3; 5; 2]

list as function

If $m: \square L$ then

$$\langle n: \square L \cdot L n \rangle m = L m$$

function \approx list

application \approx indexing

function composition \approx list composition

$$- [3; 5; 2] = [-3; -5; -2]$$

$$\text{suc } [3; 5; 2] = [4; 6; 3]$$

list as function

If $m: \square L$ then

$$\langle n: \square L \cdot L n \rangle m = L m$$

function \approx list

application \approx indexing

function composition \approx list composition

$$- [3; 5; 2] = [-3; -5; -2]$$

$$suc [3; 5; 2] = [4; 6; 3]$$

$$1 \rightarrow 21 \mid [10; 11; 12] = [10; 21; 12]$$

list as function

If $m: \square L$ then

$$\langle n: \square L \cdot L n \rangle m = L m$$

function \approx list

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function composition \approx list composition

$$- [3; 5; 2] = [-3; -5; -2]$$

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list as function

If $m: \square L$ then

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function \approx list

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list as function

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ΣL

list as function

If $m: \square L$ then

$$\langle n: \square L \cdot L n \rangle m = L m$$

function \approx list

application \approx indexing

function composition \approx list composition

$$- [3; 5; 2] = [-3; -5; -2]$$

$$\text{suc } [3; 5; 2] = [4; 6; 3]$$

$$1 \rightarrow 21 \mid [10; 11; 12] = [10; 21; 12]$$

$$\Sigma L = \Sigma \langle n: \square L \cdot L n \rangle = \Sigma n: \square L \cdot L n$$

limit

$f: \text{nat} \rightarrow \text{rat}$

limit

$f: \text{nat} \rightarrow \text{rat}$

$f_0; f_1; f_2; \dots$ is a sequence of rationals

limit

$f: \text{nat} \rightarrow \text{rat}$

$f_0; f_1; f_2; \dots$ is a sequence of rationals

$\Downarrow f$

limit

$f: \text{nat} \rightarrow \text{rat}$

$f_0; f_1; f_2; \dots$ is a sequence of rationals

$$(\uparrow m \cdot \downarrow n \cdot f(m+n)) \leq \Downarrow f \leq (\downarrow m \cdot \uparrow n \cdot f(m+n))$$

limit

$f: \text{nat} \rightarrow \text{rat}$

$f_0; f_1; f_2; \dots$ is a sequence of rationals

$$(\uparrow m \cdot \downarrow n \cdot f(m+n)) \leq \Downarrow f \leq (\downarrow m \cdot \uparrow n \cdot f(m+n))$$

$$\Downarrow n \cdot 1/(n+1)$$

limit

$f: \text{nat} \rightarrow \text{rat}$

$f_0; f_1; f_2; \dots$ is a sequence of rationals

$$(\uparrow m \cdot \downarrow n \cdot f(m+n)) \leq \Downarrow f \leq (\downarrow m \cdot \uparrow n \cdot f(m+n))$$

$$0 \leq (\Downarrow n \cdot 1/(n+1)) \leq 0$$

limit

$f: \text{nat} \rightarrow \text{rat}$

$f_0; f_1; f_2; \dots$ is a sequence of rationals

$$(\uparrow m \cdot \downarrow n \cdot f(m+n)) \leq \Downarrow f \leq (\downarrow m \cdot \uparrow n \cdot f(m+n))$$

$$0 \leq (\Downarrow n \cdot 1/(n+1)) \leq 0$$

$$\Downarrow n \cdot (-1)^n$$

limit

$f: \text{nat} \rightarrow \text{rat}$

$f_0; f_1; f_2; \dots$ is a sequence of rationals

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$$x: \text{xreal} = \exists f: \text{nat} \rightarrow \text{rat} \cdot x = \Downarrow f$$

limit

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$$\Downarrow n. 1/(n+1) = 0 = \perp$$