

- 11 (dual) One operator is the dual of another operator if it negates the result when applied to the negated operands. The zero-operand operators  $\top$  and  $\perp$  are each other's duals. If  $op_0 \neg a = \neg op_1 a$  then  $op_0$  and  $op_1$  are duals. If  $(\neg a) op_0 (\neg b) = \neg(a op_1 b)$  then  $op_0$  and  $op_1$  are duals. And so on for more operands.
- (a) Of the 4 one-operand binary operators, there is 1 pair of duals, and 2 operators that are their own duals. Find them.
- (b) Of the 16 two-operand binary operators, there are 6 pairs of duals, and 4 operators that are their own duals. Find them.
- (c) What is the dual of the three-operand operator **if then else fi** ? Express it using only the operator **if then else fi** .
- (d) The dual of a binary expression without variables is formed as follows: replace each operator with its dual, adding parentheses if necessary to maintain the precedence. Explain why the dual of a theorem is an antitheorem, and vice versa.
- (e) Let  $P$  be a binary expression without variables. From part (d) we know that every binary expression without variables of the form

$$(\text{dual of } P) = \neg P$$

is a theorem. Therefore, to find the dual of a binary expression with variables, we must replace each operator by its dual and negate each variable. For example, if  $a$  and  $b$  are binary variables, then the dual of  $a \wedge b$  is  $\neg a \vee \neg b$  . And since

$$(\text{dual of } a \wedge b) = \neg(a \wedge b)$$

we have one of the Duality Laws:

$$\neg a \vee \neg b = \neg(a \wedge b)$$

The other of the Duality Laws is obtained by equating the dual and negation of  $a \vee b$  . Obtain five laws that do not appear in this book by equating a dual with a negation.

- (f) Dual operators have value tables that are each other's vertical mirror reflections. For example, the value table for  $\wedge$  (below left) is the vertical mirror reflection of the value table for  $\vee$  (below right).

$\wedge$ :	$\top \top$ $\top \perp$ $\perp \top$ $\perp \perp$	$\top$ $\perp$ $\perp$ $\perp$
------------	--	---

$\vee$ :	$\top \top$ $\top \perp$ $\perp \top$ $\perp \perp$	$\top$ $\top$ $\top$ $\perp$
----------	--	---------------------------------------

Design symbols (you may redesign existing symbols where necessary) for the 4 one-operand and 16 two-operand binary operators according to the following criteria.

- (i) Dual operators should have symbols that are vertical mirror reflections (like  $\wedge$  and  $\vee$  ). This implies that self-dual operators have vertically symmetric symbols, and all others have vertically asymmetric symbols.
- (ii) If  $a op_0 b = b op_1 a$  then  $op_0$  and  $op_1$  should have symbols that are horizontal mirror reflections (like  $\Rightarrow$  and  $\Leftarrow$  ). This implies that symmetric operators have horizontally symmetric symbols, and all others have horizontally asymmetric symbols.

After trying the question, scroll down to the solution.

(a) Of the 4 one-operand binary operators, there is 1 pair of duals, and 2 operators that are their own duals. Find them.

§ To answer this question, I'll use the symbols I introduce in part (f). The pair of duals is:  $\bar{\top}$  (always  $\top$ ) and  $\perp$  (always  $\perp$ ). The two self-duals are:  $\bar{\perp}$  (identity) and  $\neq$  (negation).

(b) Of the 16 two-operand binary operators, there are 6 pairs of duals, and 4 operators that are their own duals. Find them.

§ To answer this question, I'll use the symbols I will introduce in part (f). The six dual pairs are:  $\bar{\Delta} \nabla$ ,  $\vee \wedge$ ,  $\geq \leq$ ,  $\leq \geq$ ,  $\Delta \bar{\nabla}$ ,  $\Delta \nabla$ . The four self-duals are:  $<$ ,  $>$ ,  $\triangleright$ ,  $\triangleleft$ .

(c) What is the dual of the three-operand operator **if then else fi**? Express it using only the operator **if then else fi**.

§ Its value table is

$\top \top \top$	$\top \top \perp$	$\top \perp \top$	$\top \perp \perp$	$\perp \top \top$	$\perp \top \perp$	$\perp \perp \top$	$\perp \perp \perp$
$\top$	$\perp$	$\top$	$\perp$	$\top$	$\top$	$\perp$	$\perp$

The dual of **if a then b else c fi** is equivalent to **if a then c else b fi**.

(d) The dual of a binary expression without variables is formed as follows: replace each operator with its dual, adding parentheses if necessary to maintain the precedence. Explain why the dual of a theorem is an antitheorem, and vice versa.

§ I will show that for every expression  $P$  without variables,  $(\text{dual of } P) = \neg P$ . I do so by induction on the structure of expression  $P$ . The two binary values give us two base cases.

$$\begin{aligned}
 & (\text{dual of } \top) && \text{use the dual-forming rules} \\
 = & \perp \\
 = & \neg \top \\
 & (\text{dual of } \perp) && \text{use the dual-forming rules} \\
 = & \top \\
 = & \neg \perp
 \end{aligned}$$

There is an induction step for each of the binary operators. Suppose (this is an induction hypothesis) that  $(\text{dual of } P) = \neg P$ . Then

$$\begin{aligned}
 & (\text{dual of } \neg P) && \text{use the dual-forming rules} \\
 = & \neg(\text{dual of } P) && \text{use the induction hypothesis} \\
 = & \neg \neg P
 \end{aligned}$$

Suppose that  $(\text{dual of } P) = \neg P$  and  $(\text{dual of } Q) = \neg Q$ . Then

$$\begin{aligned}
 & (\text{dual of } P \wedge Q) && \text{use the dual-forming rules} \\
 = & (\text{dual of } P) \vee (\text{dual of } Q) && \text{use the induction hypotheses} \\
 = & \neg P \vee \neg Q && \text{use duality law} \\
 = & \neg(P \wedge Q)
 \end{aligned}$$

And similarly for all other operators.

§(e) From  $a=b$  we get  $\neg a \neq \neg b = \neg(a=b)$   
 From **if a then b else c fi** we get **if  $\neg a$  then  $\neg c$  else  $\neg b$  fi** =  **$\neg$ if a then b else c fi**  
 From  $a=b \wedge c$  we get  $\neg a \neq \neg b \vee \neg c = \neg(a=b \wedge c)$   
 From  $a=b \vee c$  we get  $\neg a \neq \neg b \wedge \neg c = \neg(a=b \vee c)$   
 From  $a = (b \wedge c)$  we get  $\neg a \neq (\neg b \vee \neg c) = \neg(a = (b \wedge c))$

§(f)

old		$\bar{\top}$	$\mathbf{I}$	$\neg$	
new		$\bar{\top}$	$\mathbf{I}$	$\neq$	$\perp$

---

$\top$	$\top$	$\top$	$\perp$	$\perp$
$\perp$	$\top$	$\perp$	$\top$	$\perp$

old		$\forall$	$\Leftarrow$	$\Rightarrow$	$=$	$\wedge$	$\neq$									
new	$\bar{\Delta}$	$\forall$	$\geq$	$<$	$\leq$	$>$	$\underline{\Delta}$	$\wedge$	$\Delta$	$\bar{\nabla}$	$\triangleright$	$\geq$	$\triangleleft$	$\leq$	$\nabla$	$\nabla$

---

$\top \top$	$\top$	$\top$	$\top$	$\top$	$\top$	$\top$	$\top$	$\top$	$\top$	$\perp$	$\perp$	$\perp$	$\perp$	$\perp$	$\perp$	$\perp$
$\top \perp$	$\top$	$\top$	$\top$	$\top$	$\perp$	$\perp$	$\perp$	$\perp$	$\perp$	$\top$	$\top$	$\top$	$\top$	$\perp$	$\perp$	$\perp$
$\perp \top$	$\top$	$\top$	$\perp$	$\perp$	$\top$	$\top$	$\perp$	$\perp$	$\perp$	$\top$	$\top$	$\perp$	$\perp$	$\top$	$\top$	$\perp$
$\perp \perp$	$\top$	$\perp$	$\top$	$\perp$	$\top$	$\perp$	$\top$	$\perp$	$\perp$	$\top$	$\perp$	$\top$	$\perp$	$\top$	$\perp$	$\perp$