

264 (edit distance) Given two lists, write a program to find the minimum number of item insertions, item deletions, and item replacements to change one list into the other.

After trying the question, scroll down to the solution.

§ Here is the standard solution, which uses **for**-loops (Subsection 5.2.3), but **for**-loops are never necessary. We will change list  $A$  into list  $B$ . Let  $D: [(\#A+1) * [(\#B+1) * nat]]$  be an array-valued variable whose final value will be such that  $D' i j =$  (the edit distance from  $A[0;..i]$  to  $B[0;..j]$ ). So the final answer will be  $D' (\#A) (\#B)$ .

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(0)   for  $i:= 0;..\#A+1$  do  $D i 0:= i$  od.
(1)   for  $j:= 1;..\#B+1$  do  $D 0 j:= j$  od.
(2)   for  $i:= 1;..\#A+1$  do
(3)     for  $j:= 1;..\#B+1$  do
(4)        $D i j:= (D (i-1) (j-1) + \text{if } A i = B j \text{ then } 0 \text{ else } 1 \text{ fi})$ 
(5)          $\downarrow (D (i-1) j + 1)$ 
(6)          $\downarrow (D i (j-1) + 1)$  od od

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Line 0 says  $A[0;..i]$  can be changed to  $[nil]$  by  $i$  deletions.

Line 1 says  $[nil]$  can be changed to  $B[0;..j]$  by  $j$  insertions.

Lines 2 and 3 fill in the interior of  $D$ .

On line 4, if we can transform  $A[0;..i-1]$  to  $B[0;..j-1]$  in  $D (i-1) (j-1)$  steps, and if  $A i = B j$ , we have transformed  $A[0;..i]$  to  $B[0;..j]$ . But if  $A i \neq B j$  then we need to replace  $A i$  by  $B j$  which takes 1 step.

On line 5, if we can transform  $A[0;..i-1]$  to  $B[0;..j]$  in  $D (i-1) j$  steps, then we can transform  $A[0;..i]$  to  $B[0;..j]$  by deleting  $A i$ .

On line 6, if we can transform  $A[0;..i]$  to  $B[0;..j-1]$  in  $D i (j-1)$  steps, then we can transform  $A[0;..i]$  to  $B[0;..j]$  by appending  $B j$ .

The shortest way to transform  $A[0;..i]$  to  $B[0;..j]$  is the minimum of the three ways from lines 4, 5, and 6.

To prove the correctness of this solution, find invariants for the **for**-loops. Or eliminate the **for**-loops and write specifications as in Chapter 4.