

179✓ (binary exponentiation) Given natural variables  $x$  and  $y$ , write a program for  $y' = 2^x$  without using exponentiation.

§ see book Subsection 4.1.2 and scroll down for proofs. See also Subsection 5.2.3.

In the textbook on page 44 there is the solution

- 0  $y'=2^x \iff \text{if } x=0 \text{ then } x=0 \Rightarrow y'=2^x \text{ else } x>0 \Rightarrow y'=2^x \text{ fi}$
- 1  $x=0 \Rightarrow y'=2^x \iff y:=1. x:=3$
- 2  $x>0 \Rightarrow y'=2^x \iff x>0 \Rightarrow y'=2^{x-1}. y'=2 \times y$
- 3  $x>0 \Rightarrow y'=2^{x-1} \iff x'=x-1. y'=2^x$
- 4  $y'=2 \times y \iff y:=2 \times y. x:=5$
- 5  $x'=x-1 \iff x:=x-1. y:=7$

Proof of refinement 0.

$$\begin{aligned}
 & y'=2^x && \text{case creation ( } x \text{ is natural)} \\
 = & \text{if } x=0 \text{ then } x=0 \Rightarrow y'=2^x \text{ else } x>0 \Rightarrow y'=2^x \text{ fi}
 \end{aligned}$$

Proof of refinement 1.

$$\begin{aligned}
 & (x=0 \Rightarrow y'=2^x \iff y:=1. x:=3) && \text{portation} \\
 = & x=0 \wedge (y:=1. x:=3) \Rightarrow y'=2^x && \text{definition of assignment and substitution law} \\
 = & x=0 \wedge y'=1 \wedge x'=3 \Rightarrow y'=2^x && \text{context} \\
 = & x=0 \wedge y'=1 \wedge x'=3 \Rightarrow 1=2^0 && \text{arithmetic} \\
 = & x=0 \wedge y'=1 \wedge x'=3 \Rightarrow \top && \text{base} \\
 = & \top
 \end{aligned}$$

Proof of refinement 2.

$$\begin{aligned}
 & x>0 \Rightarrow y'=2^{x-1}. y'=2 \times y && \text{sequential composition} \\
 = & \exists x'', y''. (x>0 \Rightarrow y''=2^{x-1}) \wedge y'=2 \times y'' && \text{idempotent ( } x'' \text{ does not appear)} \\
 = & \exists y''. (x>0 \Rightarrow y''=2^{x-1}) \wedge y'=2 \times y'' && \text{arithmetic} \\
 = & \exists y''. (x>0 \Rightarrow y''=2^{x-1}) \wedge y''=y'/2 && \text{one-point} \\
 = & x>0 \Rightarrow y'/2 = 2^{x-1} && \text{arithmetic} \\
 = & x>0 \Rightarrow y'=2^x
 \end{aligned}$$

Proof of refinement 3.

$$\begin{aligned}
 & (x>0 \Rightarrow y'=2^{x-1} \iff x'=x-1. y'=2^x) && \text{portation} \\
 = & x>0 \wedge (x'=x-1. y'=2^x) \Rightarrow y'=2^{x-1} && \text{sequential composition} \\
 = & x>0 \wedge (\exists x'', y''. x''=x-1 \wedge y'=2^{x''}) \Rightarrow y'=2^{x-1} && \text{idempotent and one-point} \\
 = & x>0 \wedge y'=2^{x-1} \Rightarrow y'=2^{x-1} && \text{specialization} \\
 = & \top
 \end{aligned}$$

Proof of refinement 4.

$$\begin{aligned}
 & y:=2 \times y. x:=5 && \text{definition of assignment and substitution law} \\
 = & x'=5 \wedge y'=2 \times y && \text{specialization} \\
 \Rightarrow & y'=2 \times y
 \end{aligned}$$

Proof of refinement 5.

$$\begin{aligned}
 & x:=x-1. y:=7 && \text{definition of assignment and substitution law} \\
 = & x'=x-1 \wedge y'=7 && \text{specialization} \\
 \Rightarrow & x'=x-1
 \end{aligned}$$