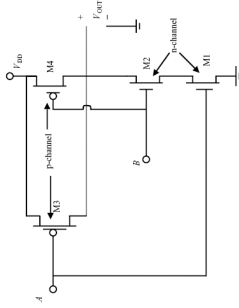


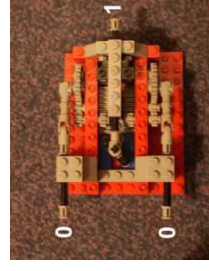
Introduction to Logic Gates

- Using transistor technology, we can create basic **logic gates** that perform boolean operations on high (5V) and low (0V) signals.
- Example: NAND gate



Intro to Logic Gates

- Digital logic gates form the basis for all implemented computing machines.



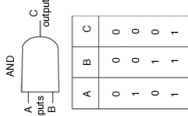
- Basic gates:
 - AND : $\cdot$, $\wedge$
 - OR : $+$, $\vee$
 - NOT : $\neg$, $-$
 - NAND
 - NOR
 - XOR : $\oplus$
 - XNOR

- Relationships between inputs and outputs are outlined in **truth tables**

Truth Tables

- AND gate

$$\rightarrow C = A \cdot B$$



A	B	C
0	0	0
1	0	0
0	1	0
1	1	1

- OR gate

$$\rightarrow C = A + B$$



A	B	C
0	0	0
1	0	1
0	1	1
1	1	1

Truth Tables (cont'd)

- Buffer



A	B
0	0
1	1

- NOT Gate

$$\rightarrow B = \bar{A}$$



A	B
0	1
1	0

Truth Tables (cont'd)

- NAND Gate



A	B	C
0	0	1
1	0	1
0	1	1
1	1	0

- NOR Gate



A	B	C
0	0	1
1	0	0
0	1	0
1	1	0

Truth Tables (cont'd)

- XOR Gate



$$\rightarrow C = A \oplus B$$

- XNOR Gate



A	B	C
0	0	1
1	0	0
0	1	0
1	1	1

Logic Operations

- Logic gates all follow the same rules as logic operators in programming languages.

AND: $1 \cdot x = x$
 $0 \cdot x = 0$
 OR: $1 + x = 1$
 $0 + x = x$
 NOT: $x + \overline{x} = 1$
 $x \cdot \overline{x} = 0$
 XOR: $x \oplus y = y \oplus x$



- Order of operations also applies to these operations, when putting together the digital circuit.

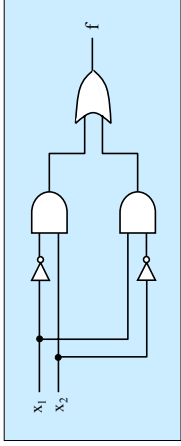
Designing a Circuit

- Task #1:** Given a logic expression, determine the equivalent gate representation.
 - Example: $f = \overline{x_1} \cdot x_2 + x_1 \cdot \overline{x_2}$
- Group terms according to **order of operations**:
 - NOT terms first
 - AND terms next
 - OR terms last
- The example gets rewritten as:

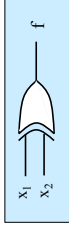
$$f = ((\overline{x_1}) \cdot x_2) + (x_1 \cdot (\overline{x_2}))$$



Circuit Example



- By evaluating the terms of the expression from the inside-out, we construct the diagram above.
- Of course, this can also be represented as a single gate, but the underlying complexity is the same.



More Circuit Design

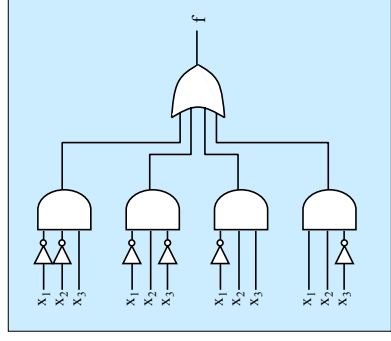
- Task #2:** Given a truth table that specifies a logic circuit's behaviour, design the equivalent circuit.
 - Example: three-input circuit
- Sum-of-product technique:**
 - Group all rows with an output of $f=1$ into a single AND term (product)
 - Combine these AND terms with a single OR gate (sum)
- Note:** All truth tables can be converted into gate form by using this technique.

x_1	x_2	x_3	f
0	0	0	0
0	0	1	1
0	1	0	1
0	1	1	1
1	0	0	0
1	0	1	0
1	1	0	1
1	1	1	0

Sum-of-Products

- A **minterm** is a product term where each of the input variables occurs exactly once.
 - The sum-of-products technique is also referred to as a **disjunction of minterms**
- Grouping together the minterms for all rows where $f=1$ yields the following expression:

$$f = \bar{x}_1 \bar{x}_2 x_3 + \bar{x}_1 x_2 \bar{x}_3 + \bar{x}_1 x_2 x_3 + x_1 x_2 \bar{x}_3$$



Reducing Minterms

- This is an ugly expression. Can we find some way to **minimize** the expression, to make it more compact?
- We can, by employing a set of binary logic rules:

$$f = \bar{x}_1 \bar{x}_2 x_3 + \bar{x}_1 x_2 \bar{x}_3 + \bar{x}_1 x_2 x_3 + x_1 x_2 \bar{x}_3$$

Rule Name	Algebraic Identity
Commutative	$x + y = y + x$ $xy = yx$
Associative	$(x + y) + z = x + (y + z)$ $(xy)z = x(yz)$
Distributive	$x + yz = (x + y)(x + z)$ $x(y + z) = xy + xz$
Idempotent	$x + x = x$ $xx = x$
Involution	$\bar{\bar{x}} = x$
Complement	$x + \bar{x} = 1$ $x\bar{x} = 0$
de Morgan	$\overline{x + y} = \bar{x} \cdot \bar{y}$ $\overline{xy} = \bar{x} + \bar{y}$

Simplification Example

- Reduce the following sum of products:

$$f = \bar{x}_1 \bar{x}_2 x_3 + \bar{x}_1 x_2 x_3 + x_1 \bar{x}_2 x_3 + x_1 x_2 \bar{x}_3 + x_1 x_2 x_3$$

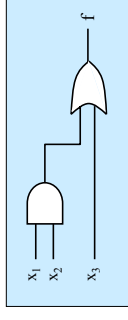
- Steps:
 - Distributive: $f = \bar{x}_1 x_3 (\bar{x}_2 + x_2) + x_1 \bar{x}_2 x_3 + x_1 x_2 \bar{x}_3 + x_1 x_2 x_3$
 - Complement: $f = \bar{x}_1 x_3 \cdot 1 + x_1 \bar{x}_2 x_3 + x_1 x_2 \bar{x}_3 + x_1 x_2 x_3$
 - Identity: $f = \bar{x}_1 x_3 + x_1 \bar{x}_2 x_3 + x_1 x_2 \bar{x}_3 + x_1 x_2 x_3$
 - Idempotent: $f = \bar{x}_1 x_3 + x_1 \bar{x}_2 x_3 + x_1 x_2 \bar{x}_3 + (x_1 x_2 \bar{x}_3 + x_1 x_2 x_3)$
 - Associative: $f = \bar{x}_1 x_3 + (x_1 \bar{x}_2 x_3 + x_1 x_2 x_3) + (x_1 x_2 \bar{x}_3 + x_1 x_2 x_3)$
 - Distributive: $f = \bar{x}_1 x_3 + x_1 x_3 (\bar{x}_2 + x_2) + x_1 x_2 (\bar{x}_3 + x_3)$
 - Complement: $f = \bar{x}_1 x_3 + x_1 x_3 \cdot 1 + x_1 x_2 \cdot 1$
 - Identity: $f = \bar{x}_1 x_3 + x_1 x_3 + x_1 x_2$

Simplification Example

- Steps (cont'd):
 - (from previous slide): $f = \bar{x}_1 x_3 + x_1 x_3 + x_1 x_2$
 - Distributive: $f = x_3 (\bar{x}_1 + x_1) + x_1 x_2$
 - Complement: $f = x_3 \cdot 1 + x_1 x_2$
 - Identity: $f = x_3 + x_1 x_2$

- Solution is:

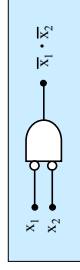
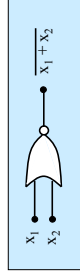
$$f = x_3 + x_1 x_2$$



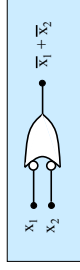
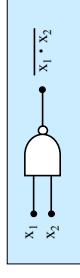
More de Morgan

- Implications of de Morgan's Law:

$$\overline{x + y} = \bar{x} \cdot \bar{y}$$



$$\overline{x \cdot y} = \bar{x} + \bar{y}$$



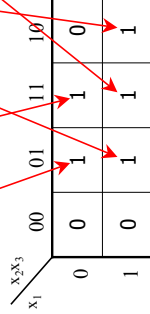
Karnaugh Maps

- A **Karnaugh map** is another way of representing a circuit's truth table
 - Presented as a two-dimensional grid of 2^n squares.
 - Each axis represents the different input values to the circuit, and the contents of each square represents the output value for the intersecting input values.
 - Horizontally and vertically adjacent squares only differ by a single input variable
 - Number of rows and columns must be a power of 2
 - Note: row and column labels must also differ by a single digit.
 - Simplified circuit is found by circling groups of adjacent 1's on the grid.
 - Result: The minimal expression for the circuit.

Karnaugh Maps

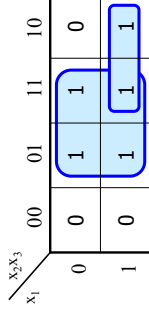
- Example:

$$f = \bar{x}_1 \bar{x}_2 x_3 + \bar{x}_1 x_2 x_3 + x_1 \bar{x}_2 x_3 + x_1 x_2 \bar{x}_3 + x_1 x_2 x_3$$



Karnaugh Maps

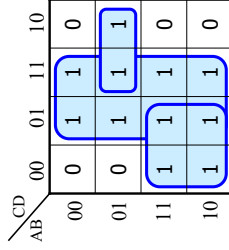
- Next step: circle blocks of 1's
 - Cannot contain any 0's in the block
 - Height and width of block must be a power of 2
 - Blocks are allowed to overlap



- Circled blocks correspond to x_3 and $x_1 x_2$

Karnaugh Example

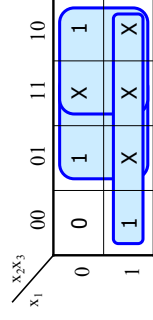
- Task: Given the truth table on the right, determine the simplest equivalent gate arrangement.



$$X = A\bar{C} + D + \bar{A}BC$$

"Don't Care" Conditions

- Sometimes certain outputs are not defined
 - Example: A three-input circuit where the output is 1 if any of the inputs is 1, and 0 if all of the inputs are 0.
 - The output can be anything in the case where two or more inputs are 1, since we don't consider those cases.
 - In those cases, label the output as "don't care" (or "X").
- "Don't care" conditions are useful because they can be treated as either 1 or 0, depending on which makes things easier.
 - In real life, you usually have to set it to some sensible value.



"Don't Care" Example

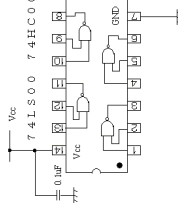
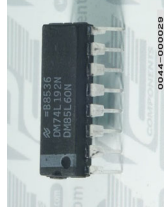
- Adding milk to beverages:
 - If a patron orders coffee (C), add milk. If the patron orders tea (T), only add milk if lemon (L) hasn't already been added.

		CT			
L		00	01	11	10
	0	X	1	X	1
	1	X	0	X	1

- The "don't care" conditions are sometimes written as "d" (as in the textbook, for example). In industry though, "X" is more commonly used, as in the phrase "X-propagation".

The Importance of NAND

- NAND gates are considered to be the "universal" gate, because any other gate can be synthesized using NAND.
- In fact, most gates are implemented in solid-state TTL chips (Transistor-Transistor Logic)
 - e.g. 74LS00 integrated circuit (IC)

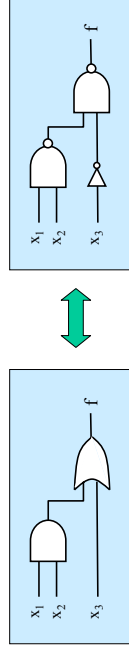


Getting to NAND

- How do we create a NAND-based circuit, if our approach so far has been to find sum-of-products?
 - Answer: By using de Morgan's rule!
- According to de Morgan (see slide 15),



- Therefore, the following are equivalent as well:



Neither...NOR

- NAND and NOR gates are the cheapest to make, and the most commonly found in IC chips
- Typical process for creating a NAND-based circuit is:
 - represent truth table in Karnaugh map
 - isolate smallest terms that produce high output
 - produce sum-of-product model with AND & OR gates
 - convert model to NAND representation
- What about creating an equivalent circuit out of NOR gates? Is that possible?
 - Requires obtaining **maxterms** of the truth table, and producing a **product-of-sums** model

Maxterms

- A **maxterm** is a sum of input values where every input value occurs exactly once.
- When expressing a truth table, each maxterm represents a row where the output is 0, such that the maxterm will give a true value in all cases except for that row.
 - Maxterms for example table:
 - $(x_1 + x_2 + x_3)$
 - $(\bar{x}_1 + x_2 + x_3)$
 - $(\bar{x}_1 + x_2 + \bar{x}_3)$
 - $(\bar{x}_1 + \bar{x}_2 + \bar{x}_3)$

x_1	x_2	x_3	f
0	0	0	0
0	0	1	1
0	1	0	1
0	1	1	1
1	0	0	0
1	0	1	0
1	1	0	1
1	1	1	0

POS (Product of Sums)

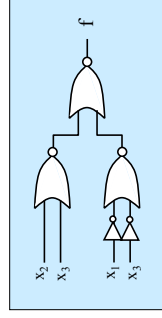
- Instead of creating a disjunction of cases where the output of the circuit is 1, the **product-of-sums** technique creates a conjunction of the cases where the output is 0.
 - From previous example:

$(x_1 + x_2 + x_3) \cdot (\bar{x}_1 + x_2 + x_3) \cdot (\bar{x}_1 + x_2 + \bar{x}_3) \cdot (\bar{x}_1 + \bar{x}_2 + \bar{x}_3)$
- These equations can be reduced using the same techniques used on minterms.
 - boolean logic rules
 - Karnaugh maps

NOR Circuit

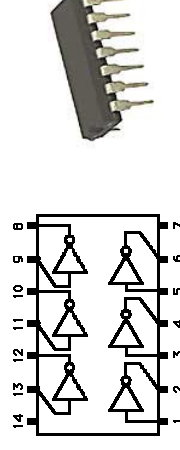
- A product-of-sum circuit will result in several OR gates, united by a single AND gate.
- Through de Morgan's rule, this converts to a circuit made up entirely of NOR gates.
- Design decision: of the maxterm and minterm representations, which has the fewest reduced terms?

$x_2 x_3$	00	01	11	10
0	0	1	1	1
1	0	0	0	1



Other TTL Logic

- In addition to NAND circuits, inverter circuits are commonly used as well (e.g. 74LS04).



- When using these chips, make sure you have the pins and the orientation correct. For example, applying the voltage source to the ground pin can have some very unpleasant results.