Codifying Probabilities with The Last Full Measure

Albert Lai

November 1, 2023

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Introduction

Motivating Example

Laplace's Rule of Succession

Let $n, k \in \mathbb{N}$, $0 \le k \le n$, be given/fixed.

- 1. pick *r* from Uniform(0, 1)
- 2. pick c from Binomial(n, r)
- 3. on condition / restrict to c = k

Expected value of r = (k + 1)/(n + 2)

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- Hybrid continuous-discrete.
- Randomly choose a binomial distribution. ••

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(Finishing what 3Blue1Brown started: probabilities of probabilities part 1, part 2.)

Programming Probabilities

Random Number Generators

RNG*X* is set/type of programs that output a random *X* element (probability depends on program)

Uniform : $\mathbb{R} \times \mathbb{R} \to \text{RNG} \mathbb{R}$ Uniform(a, b) picks real number from [a, b] evenly.

Binomial : $\mathbb{N} \times [0, 1] \rightarrow \text{RNG} \mathbb{N}$

Binomial(n, r) tosses coin *n* times, head probability *r*, count heads.

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Binomial : $\mathbb{N} \times [0, 1] \to \text{RNG } \mathbb{N}$ Binomial(n, r) tosses coin *n* times, head probability *r*, count heads.

Unit_{*X*} : $X \rightarrow \text{RNG} X$ Unit_{*X*}(*x*) Always outputs *x*. "Rare desert of determinism in vast oasis of randomization." I omit subscript *X* if inferrable.

Chaining

Chaining operator "bind", "then", "flat-map" \gg = : RNG *X* × (*X* → RNG *Y*) → RNG *Y*

 $g \gg = k$ passes output of g to parametrized RNG k

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Example:

- 1. pick *r* from Uniform(0, 1)
- **2**. pick *c* from Binomial(n, r)
- **3**. output (*r*, *c*)

Uniform(0, 1) $\gg = (\lambda r \cdot \text{Binomial}(n, r) \gg = (\lambda c \cdot \text{Unit}(r, c))$ or

 $\text{Uniform}(0,1) \gg = \lambda r \cdot \text{Binomial}(n,r) \gg = \lambda c \cdot \text{Unit}(r,c)$

Restriction/Conditional

Restriction operator |: RNG $X \times (X \rightarrow \mathbb{B}) \rightarrow \text{RNG } X$ ($\mathbb{B} = \{\text{false, true}\}$)

 $g \mid pred = restrict g$ to when *pred* is true.

(Rejection sampling: Keep retrying g until output satisfies pred.)

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Example:

- 1. pick *x* from Uniform(0, 1)
- 2. on condition x > 0.3
- **3**. (output *x*)

Uniform(0, 1) | ($\lambda x \cdot x > 0.3$)

Express as Program

Define

$$g = \text{Uniform}(0, 1) \gg = \lambda r \cdot \text{Binomial}(n, r) \gg = \lambda c \cdot \text{Unit}(r, c)$$
$$g' = g \mid (\lambda(r, c) \cdot c = k)$$
$$\mu = g' \gg = \lambda(r, c) \cdot \text{Unit}(r)$$

Then μ generates an r as prescribed.

Re-run many times to approximate distribution and expected value. (Actual Haskell code modulo syntax.)

Re-read Program as Probability Measure

There is also a measure-theory reading of

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Then μ is the probability distribution/measure of the *r* in question. Can find expected value and pdf.

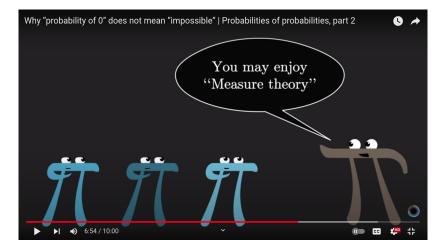
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"Codifying probability with the last full measure."



Probability Measures

Measure Theory: Motivation 1/2

Integration by Lateral Thinking. Literally.

Riemann integral: Pixelate the *x*-axis.

 $\sum_{i} f(x_i) \times \text{length}[x_i, x_i + \epsilon)$

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Lebesgue integral: Pixelate the *y*-axis.

$$\sum_i y_i \times \operatorname{length} \left(f^{-1}[y_i, y_i + \epsilon) \right)$$

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$$\sum_{i} y_i \times \operatorname{length} \left(f^{-1}[y_i, y_i + \epsilon) \right)$$

Need "length" for fairly general subsets. Require:

 $length(\bigcup_i A_i) = \sum_i length(A_i)$ (countable disjoint union)

Turns out problematic for all subsets. Settle for large enough family closed under complement, countable union.

Measure Theory: Motivation 2/2

Probability Theory

Sample space Ω .

```
Pr : subsets of \Omega \rightarrow [0, 1]
```

Require:

$$\Pr(\bigcup_i A_i) = \sum_i \Pr(A_i)$$
 (countable disjoint union)

Hmm déjà vu...

Measure Theory: Motivation 2/2

Probability Theory

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Require:

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Moreover, expected value of $f : \Omega \to \mathbb{R}$ = use \Pr for "length" in Lebesgue integration!

$$\sum_i y_i \times \Pr\left(f^{-1}[y_i, y_i + \epsilon)\right)$$

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Measure Theory: Motivation 2/2

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Riemann may not work: Ω may not even have "intervals".

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Data of a measurable space:

- Set of points X.
- σ -algebra \mathfrak{M} or \mathfrak{M}_X , $\subseteq \mathfrak{P}(X)$, closed under:
 - ▶ owning X, owning Ø
 - complement, countable union
 - (corollary: also subtraction, countable intersection)

Members "measurable subsets".

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- X = ℝ, 𝔐 Borel algebra: smallest σ-algebra owning open sets.
- ► X = ℝ, 𝔐 Lebesgue-measurable subsets (appendix). Larger than Borel algebra. Default.

Measures; Probability Measures

[Positive] Measure μ over measurable space (X, \mathfrak{M}):

- $\blacktriangleright \mu: \mathfrak{M} \to [0,\infty]$
- $\mu(\bigcup_i A_i) = \sum_i \mu(A_i)$ for countable disjoint union "countably additive"
- (corollary: $\mu(\emptyset) = 0$)
- (corollary: if $A \subseteq B$ then $\mu(A) \le \mu(B)$)

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Probability measure a.k.a. distribution: Furthermore:

$$\mu(X) = 1$$

• (corollary: $\mu : \mathfrak{M} \to [0, 1]$)

Define ΠX = set of probability measures over (X, \mathfrak{M}_X) .

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Discrete Measures

Counting measure: *X* countable set, $\mathfrak{M} = \wp(X)$ # : $\wp(X) \to \mathbb{N} \cup \{\infty\} \subseteq [0, \infty]$

$$#(A) = \begin{cases} |A| & A \text{ finite} \\ \infty & A \text{ infinite} \end{cases}$$

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Binomial distribution: $X = \mathbb{N}$ (for simplicity), $\mathfrak{M} = \wp(X)$ Binomial : $\mathbb{N} \times [0, 1] \to \Pi \mathbb{N}$

Binomial(n, r)(A) =
$$\sum_{\substack{i \in A \\ 0 \le i \le n}} {\binom{n}{i}} r^i (1-r)^{n-i}$$

Continuous Measures

Lebesgue measure: $X = \mathbb{R}$, \mathfrak{M} Lebegue-measurable subsets. $m : \mathfrak{M} \to [0, \infty]$ m([a, b]) = b - a, same for (a, b] etc. Full defn in appendix. Default.

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Uniform distribution: Lebesgue measure with rescaling: Uniform : $\mathbb{R}\times\mathbb{R}\to\Pi\,\mathbb{R}$

Uniform
$$(a, b)(A) = \frac{1}{b-a}m(A \cap [a, b])$$

Continuous Measures

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Other continuous distributions mentioned later after integration.

The Unit Measure

Unit_X : X
$$\rightarrow \Pi X$$

Unit_X(x)(A) = $\chi_A(x) = \begin{cases} 0 & x \notin A \\ 1 & x \in A \end{cases}$
(χ_A characteristic function of set A)

Deterministic corner case.

X can be discrete or continuous or any measurable space.

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Deterministic corner case.

X can be discrete or continuous or any measurable space.

In the continuous case, no probability density function.



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Measurable Functions

We will only integrate functions that satisfy:

 $f: X \to [-\infty, \infty]$ measurable function iff any of:

• for all open
$$B, f^{-1}(B) \in \mathfrak{M}_X$$

• for all
$$y \in \mathbb{R}$$
, $\{x \mid f(x) > y\} \in \mathfrak{M}_X$

Motivation: Need $\mu(f^{-1}(y, y + \epsilon))$, makes sense for \mathfrak{M}_X only.

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or ≥, or <, or ≤</p>

Motivation: Need $\mu(f^{-1}(y, y + \epsilon))$, makes sense for \mathfrak{M}_X only.

Easy: $f : \mathbb{N} \to [-\infty, \infty]$ is measurable using $\mathfrak{M} = \wp(\mathbb{N})$.

Theorem: Piecewise continuous $f : \mathbb{R} \to \mathbb{R}$ is measurable using Lebesgue-measurable subsets.

Lebesgue Integration: Idea

Notation: $\int f d\mu$ integrates *f* over all of *X* using measure μ .

Assume $f \ge 0$ for now. Approximate from below by simple functions (finite range).

Example: $X = [-4, 4], f(x) = x^2$, an approximation is

$$s(x) = 4 \cdot \chi_{(-3,-2] \cup [2,3)}(x) + 9 \cdot \chi_{[-4,-3] \cup [3,4]}(x)$$

"Clearly"

$$\int s \, d\mu = 4 \cdot \mu((-3, -2] \cup [2, 3)) + 9 \cdot \mu([-4, -3] \cup [3, 4])$$

Take supremum over all approximations.

For general f, split into positive and negative parts, both treatable as above.

Lebesgue Integration

Simple functions (finite range):

$$\int \sum_{i=1}^{n} a_i \,\chi_{E_i} \,d\mu = \sum_{i=1}^{n} a_i \,\mu(E_i) \quad (a_i \in \mathbb{R}, \, E_i \in \mathfrak{M})$$

Extend to non-negative $f: X \to [0, \infty]$

$$\int f \, d\mu = \sup\{\int s \, d\mu \mid s \text{ simple, } 0 \le s \le f\}$$

Extend to full range $f: X \to [-\infty, \infty]$

$$\int f \, d\mu = \left(\int \max(0, f) \, d\mu\right) - \left(\int -\min(f, 0) \, d\mu\right)$$

assuming not $\infty - \infty$.

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Lebesgue Integration

Notation: $\int_A f d\mu$ is over arbitrary $A \in \mathfrak{M}$, instead of all *X*.

Two equivalent treatments:

- ▶ Revise the definitions, change $\mu(E_i)$ to $\mu(E_i \cap A)$.
- ► Just use $\int f \chi_A \ d\mu$

Discrete Integration

X countable set, $\mathfrak{M} = \wp(X)$, counting measure:

$$\int f \, d\# = \sum_{i \in X} f(i)$$

(assume $f \ge 0$ or else absolute convergence or other conditions)

Discrete Integration

X countable set, $\mathfrak{M} = \wp(X)$, counting measure:

$$\int f \, d\# = \sum_{i \in X} f(i)$$

(assume $f \ge 0$ or else absolute convergence or other conditions) Binomial expressible as integral d#:

Binomial(n, r)(A) =
$$\int_{A} \lambda i \cdot {\binom{n}{i}} r^{i} (1-r)^{n-i} d\#$$

Hence $\lambda i \cdot {n \choose i} r^i (1-r)^{n-i}$ is probability mass function.

$$\int f d(\operatorname{Binomial}(n,r)) = \int \lambda i \cdot f(i) {n \choose i} r^i (1-r)^{n-i} d\#$$

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Continuous Integration (Sorry!)

 $\int f \, dm$ default for $X = \mathbb{R}$ or subspace. Riemann toolbox reusable because:

Theorem: For Riemann-integrable $f : [a, b] \rightarrow \mathbb{R}$,

$$\int_{[a,b]} f \, dm = \operatorname{Riemann} \int_{a}^{b} f(x) \, dx$$

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Uniform distribution: Reuse *m* with rescaling:

$$\int f d(\text{Uniform}(a, b)) = \frac{1}{b-a} \int_{[a,b]} f dm$$

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Most continuous distributions have probability density functions:

$$\mu(A) = \int_{A} \operatorname{pdf} dm$$
$$\int f \, d\mu = \int f \operatorname{pdf} dm$$

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Unit Integration

Sifting property:

$$\int f d(\operatorname{Unit}(x)) = f(x)$$

Integral well defined (and beautiful), even in continuous case, even if no pdf. The power of measure theory!

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Chaining Probability Measures

Chaining operator "bind", "then", "flat-map" $\gg = : \Pi X \times (X \to \Pi Y) \to \Pi Y$

Intuition for $(\mu \gg = k)(B)$: For each $x \in X$, $k(x) \in \Pi Y$, k(x)(B) is probability of *B*. Average over *X* according to μ . (Total probability).

Likewise for expected values.

$$(\mu \gg = k)(B) = \int \lambda x \cdot k(x)(B) \, d\mu$$
$$\int f \, d(\mu \gg = k) = \int \left(\lambda x \cdot \int f \, d(k(x))\right) \, d\mu$$

(Require $k : X \to \Pi Y$ measurable function. Boils down to $(\lambda x \cdot k(x)(B)) : X \to [0, 1]$ measurable function for all $B \in \mathfrak{M}_Y$.)

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Interlude

Happy Belated Halloween! @

 $(\Pi, Unit, \gg=)$ is a programmer-friendly version of...

The Giry monad! Due to Michèle Giry. More in appendix.

Restricting Probability Measures

Restriction operator

 $|:\Pi X\times (X\to \mathbb{B})\to\Pi X$

Conditional probability.

$$(\mu \mid pred)(A) = \mu(A \cap B)/\mu(B)$$
$$\int f d(\mu \mid pred) = \frac{1}{\mu(B)} \int f \chi_B \ d\mu$$

where $B = \{x \mid pred(x)\}$, assuming $\in \mathfrak{M}$.

(One may say: pred is a measurable predicate.)

Re-read Program as Probability Measure

Define

$$g = \text{Uniform}(0, 1) \gg = \lambda r \cdot \text{Binomial}(n, r) \gg = \lambda c \cdot \text{Unit}(r, c)$$
$$g' = g \mid (\lambda(r, c) \cdot c = k)$$
$$\mu = g' \gg = \lambda(r, c) \cdot \text{Unit}(r)$$

Then μ is the probability measure of the *r* in question.

expected =
$$\int \lambda r \cdot r \, d\mu$$

= $(n+1) \int_0^1 r {n \choose k} r^k (1-r)^{n-k} \, dr$ appendix
= $\frac{k+1}{n+2}$
pdf $(r) = (n+1) {n \choose k} r^k (1-r)^{n-k}$

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Conclusion

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Formal constructs for expressing probabilistic models.

Program interpretation gives executable samplers.

Meaure-theory interpretation gives resulting distributions.

Works the same way for discrete, continuous, hybrid, nested.

Bibliography

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Appendix

Lebesgue Measure And Measurable Subsets

The following is well-defined for all $A \in \wp(\mathbb{R})$; denoted $m^*(A)$:

- ▶ let $\{I_n\}$ be sequence of disjoint open intervals, $\bigcup_n I_n \supseteq A$
- $\sum_{n} \text{length}(I_n)$ uncontroversial
- take infinum over all possibilities

Only issue: Not countably additive for some subsets.

Carathéodory's solution: Restrict to

 $\mathfrak{M} = \{A \mid \forall B \cdot m^*(B) = m^*(B \cap A) + m^*(B - A)\}$

Adopted for Lebesgue-measurable subsets.

Then define Lebesgue measure $m = m^*|_{\mathfrak{M}}$

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The Giry Monad 1/2

Prologue: Generalize "measurable functions" from $X \to [-\infty, \infty]$ to $X \to Y$, with *Y* any measurable space:

 $h: X \to Y$ measurable iff $\forall B \in \mathfrak{M}_Y \cdot h^{-1}(B) \in \mathfrak{M}_X$

 Π as endofunctor (on category of measurable spaces):

Object map: ΠX = set of probability measures over (X, \mathfrak{M}_X) . σ -algebra: Smallest s.t. for all $A \in \mathfrak{M}_X$, $(\lambda \mu \cdot \mu(A)) : \Pi X \to [0, 1]$ measurable function.

Morphism map: For measurable $h : X \to Y$, $\prod h : \prod X \to \prod Y$

$$(\Pi h)(\mu)(B) = \mu(h^{-1}(B))$$
$$\int f d((\Pi h)(\mu)) = \int f \circ h \, d\mu$$

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The Giry Monad 2/2

 Π as monad:

 $\operatorname{Unit}_X : X \to \Pi X$ as defined earlier.

Multiplication $\operatorname{Flat}_X : \Pi(\Pi X) \to \Pi X$

$$\operatorname{Flat}(\mu)(A) = \int \lambda v \cdot v(A) \, d\mu$$
$$\int f \, d(\operatorname{Flat}(\mu)) = \int \left(\lambda v \cdot \int f \, dv\right) \, d\mu$$

Flattens "probability of probabilities" to total probability.

Then define $\mu \gg = k = \text{Flat}((\prod k)(\mu))$.

Detailed Calculation 1/2

Let
$$B = \{(r, c) \mid c = k\} = [0, 1] \times \{k\}.$$

$$\int f \, dg = \int \left(\lambda r \cdot \int \left(\lambda c \cdot \int f \, d(\operatorname{Unit}(r, c))\right) \, d(\operatorname{Binomial}(n, r))\right) \\ \quad d(\operatorname{Uniform}(0, 1)) \\ = \int_0^1 \sum_{c=0}^n f(r, c) \binom{n}{c} r^c (1 - r)^{n-c} \, dr \\ g(B) = \int \chi_B \, dg \\ = \int_0^1 \binom{n}{k} r^k (1 - r)^{n-k} \, dr \\ = \binom{n}{k} \frac{k!(n-k)!}{(n+1)!} \qquad (\text{Beta functions}) \\ = \frac{1}{n+1}$$

Detailed Calculation 2/2

$$\int \lambda r \cdot r \, d\mu = \int \left(\lambda(r, c) \cdot \int \lambda r \cdot r \, d(\operatorname{Unit}(r)) \right) \, dg'$$

$$= \int \left(\lambda(r, c) \cdot r \right) \, dg'$$

$$= \frac{1}{g(B)} \int \left(\lambda(r, c) \cdot r \, \chi_B(r, c) \right) \, dg$$

$$= \frac{1}{g(B)} \int_0^1 r \binom{n}{k} r^k (1 - r)^{n-k} \, dr$$

$$= (n+1) \binom{n}{k} \frac{(k+1)!(n-k)!}{(n+2)!} \quad \text{(Beta functions)}$$

$$= \frac{k+1}{n+2}$$

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